

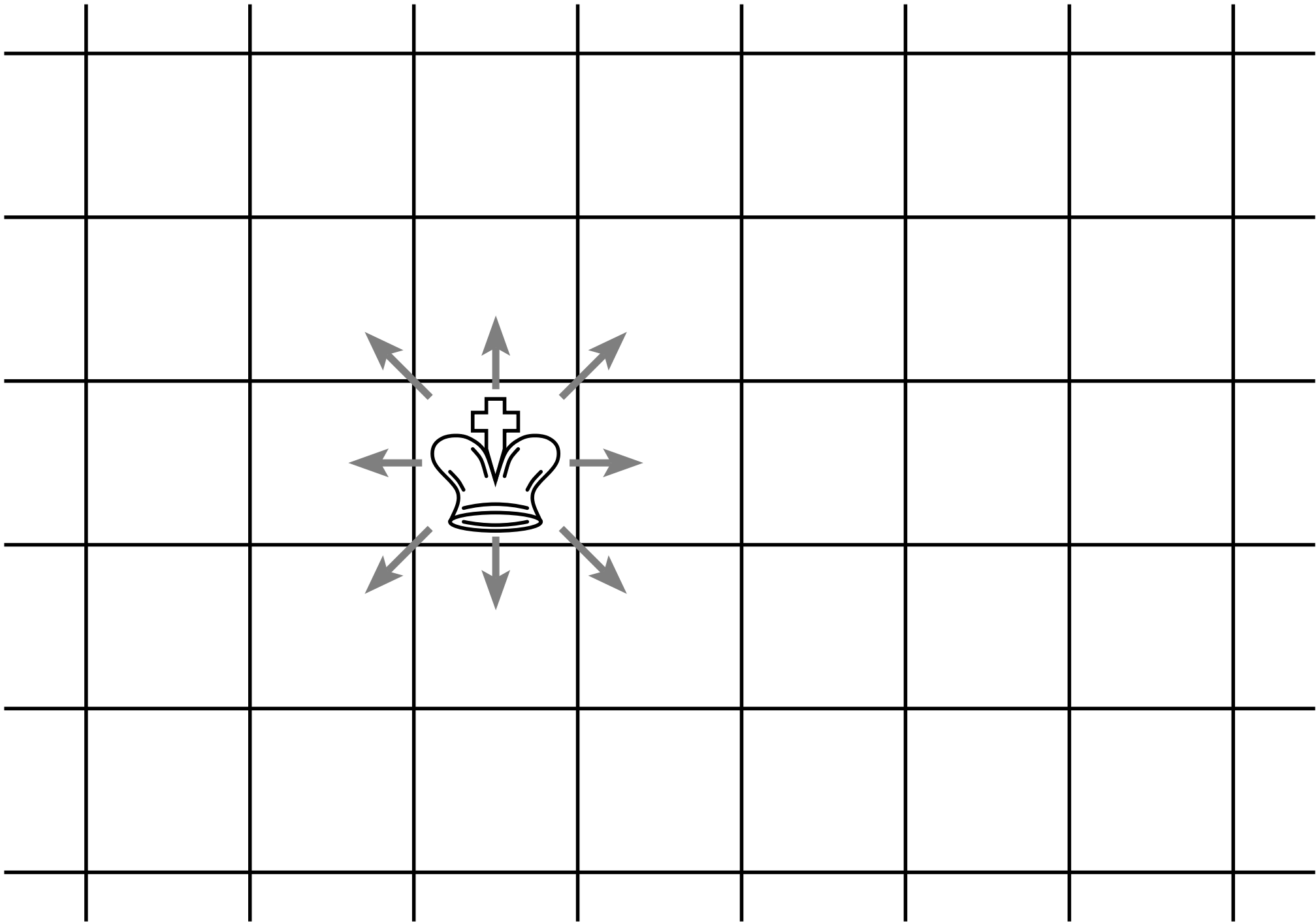
Angel, Devil, and King

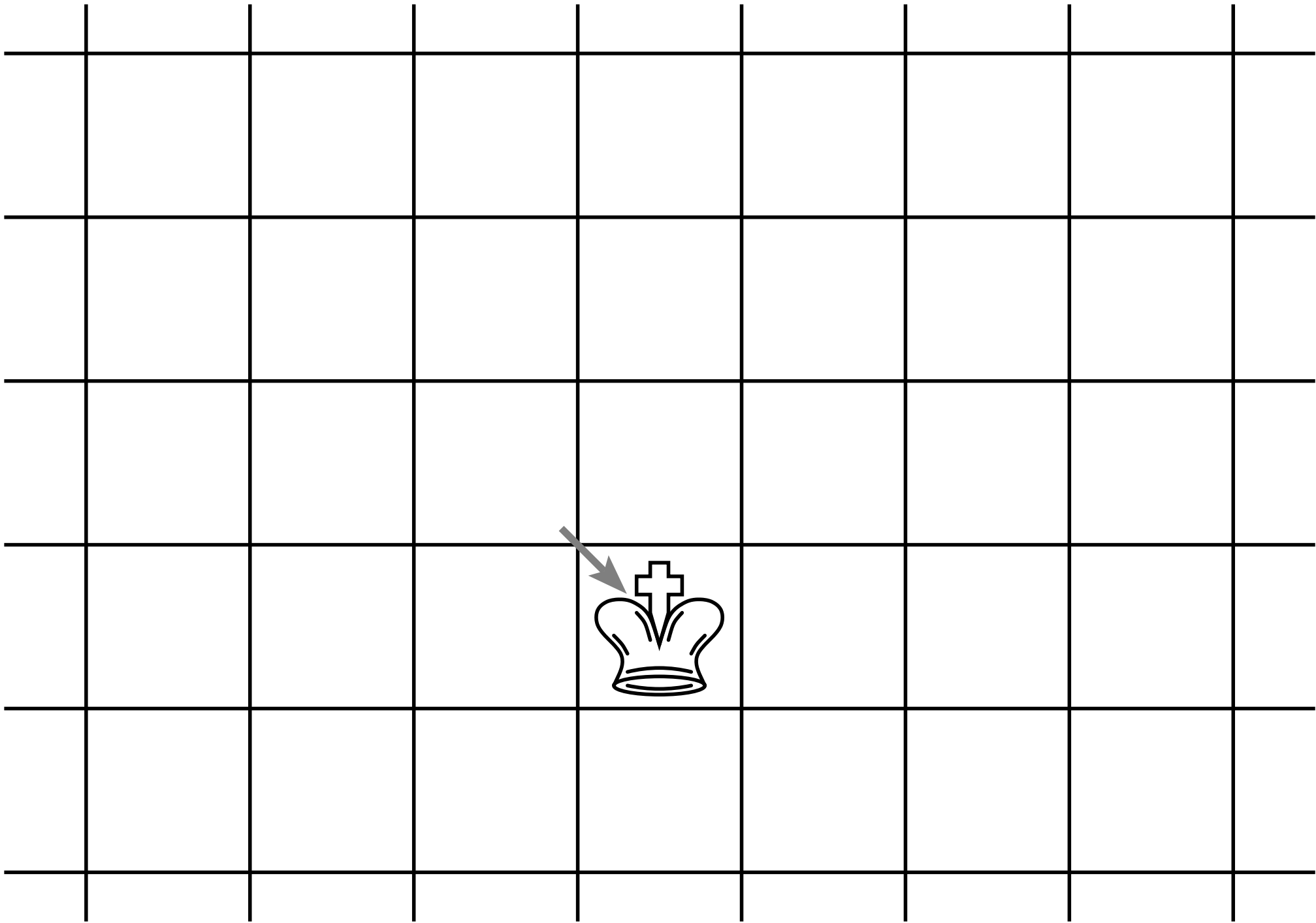
Martin Kutz

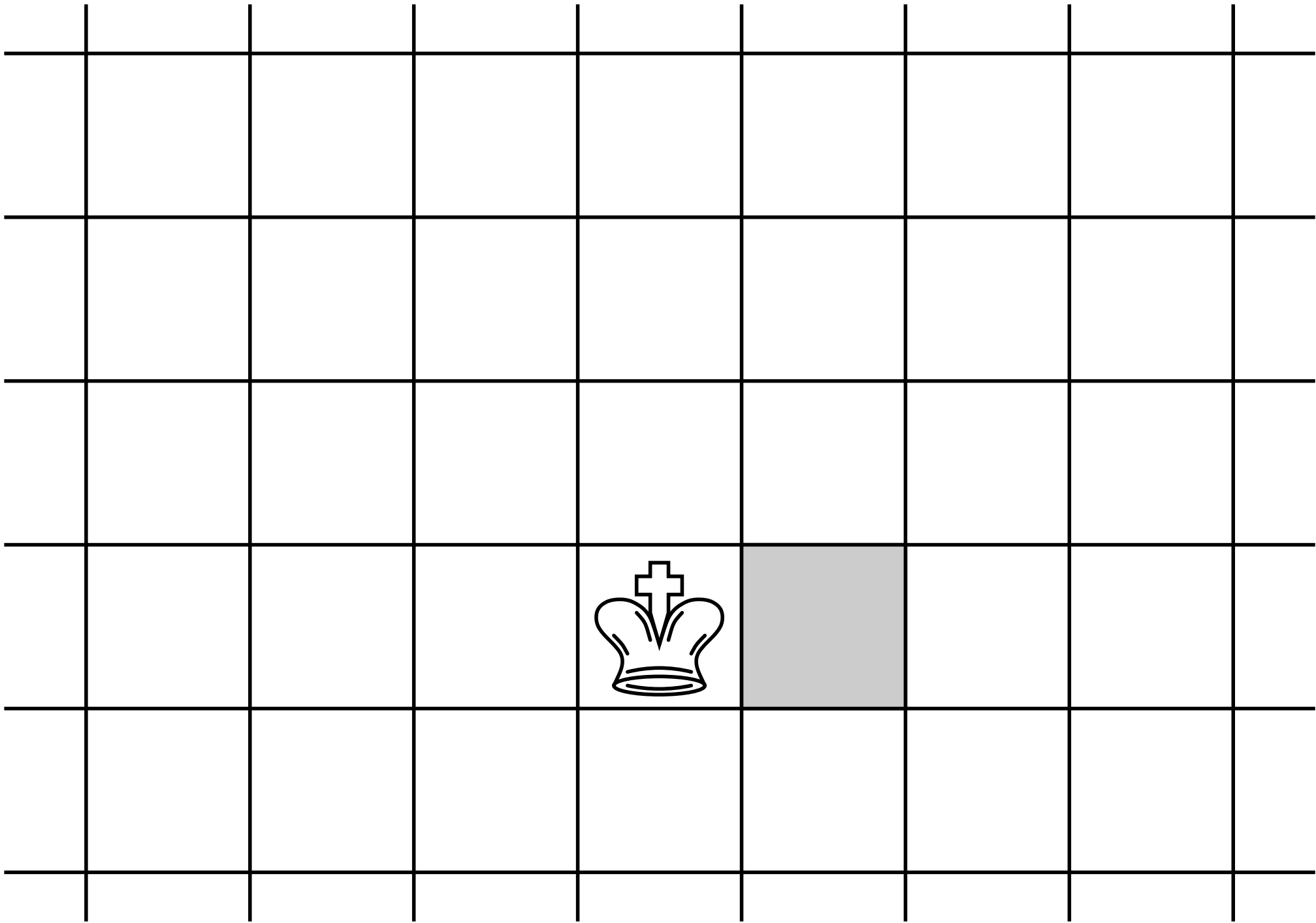
Max-Planck Institut für Informatik, Saarbrücken, Germany

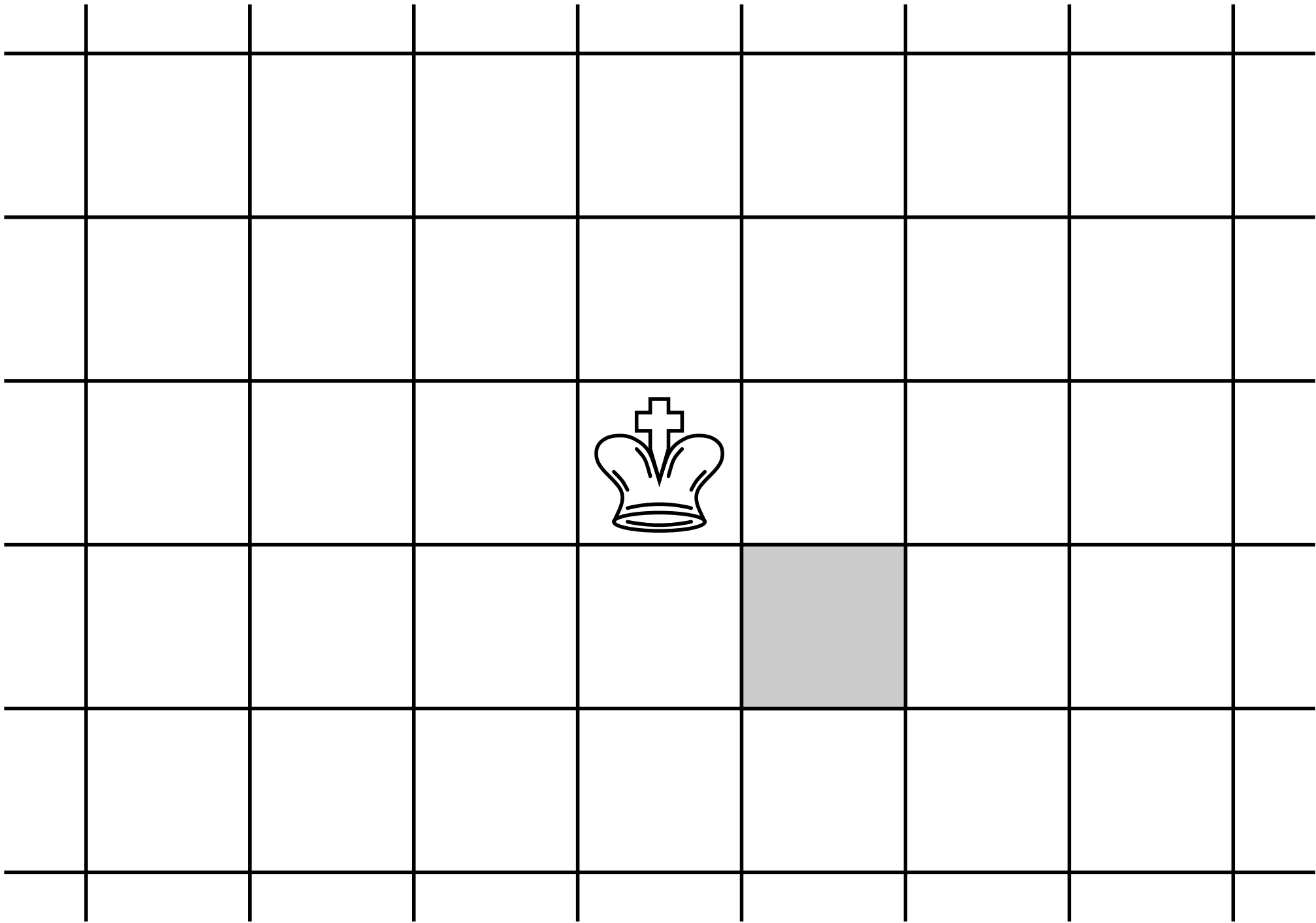
Attila Pór

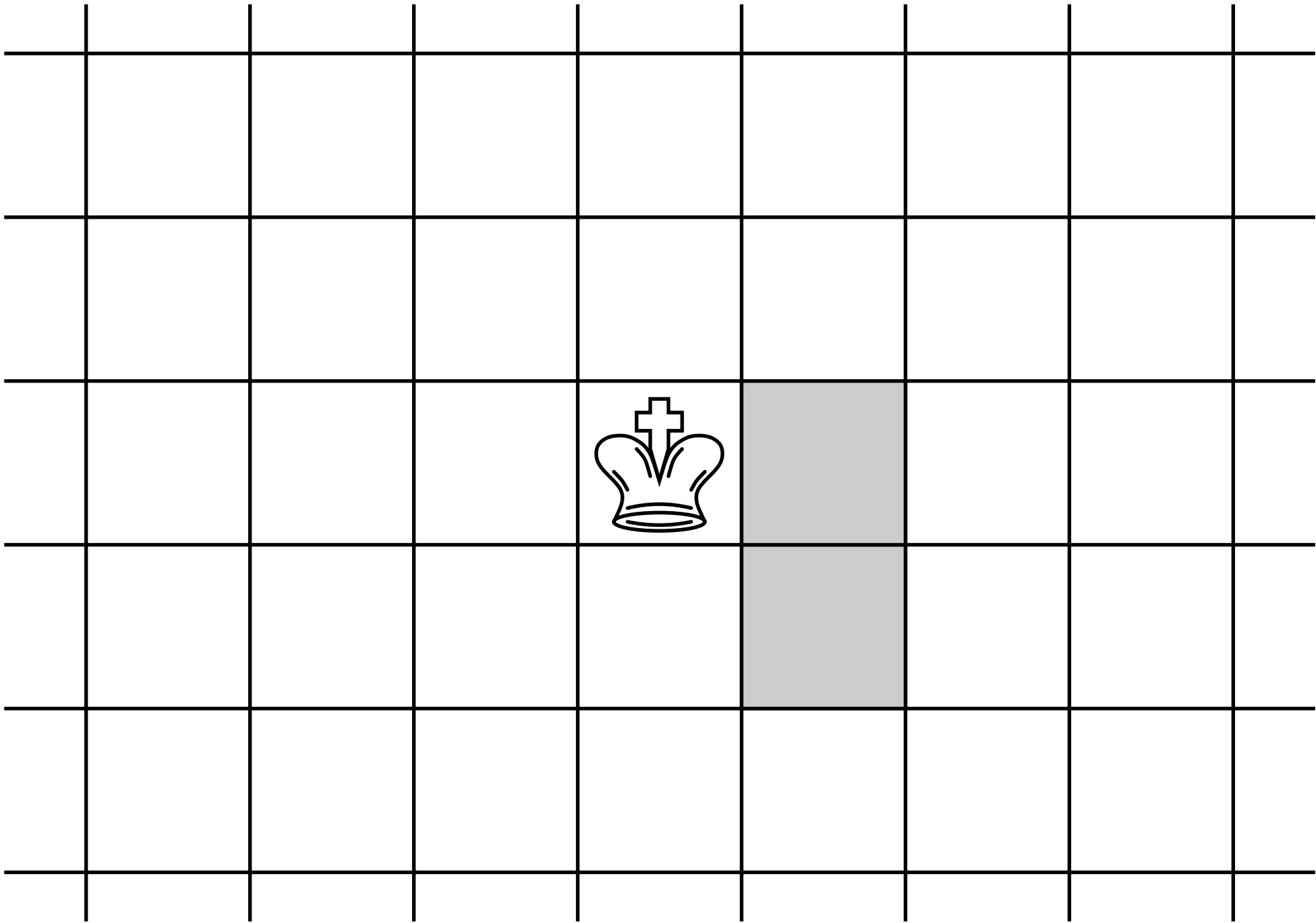
CASE Western Reserve University, Cleveland, USA

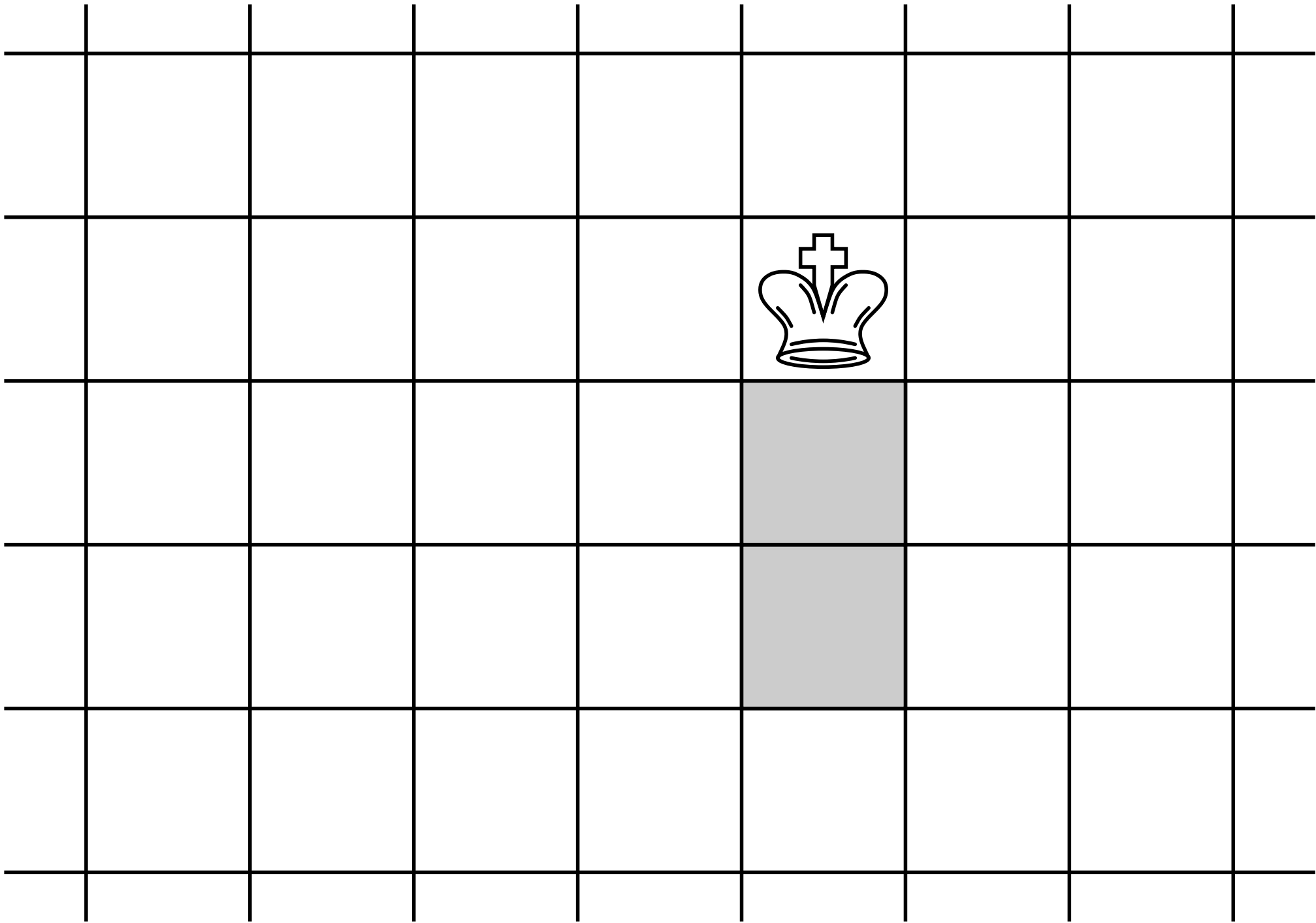


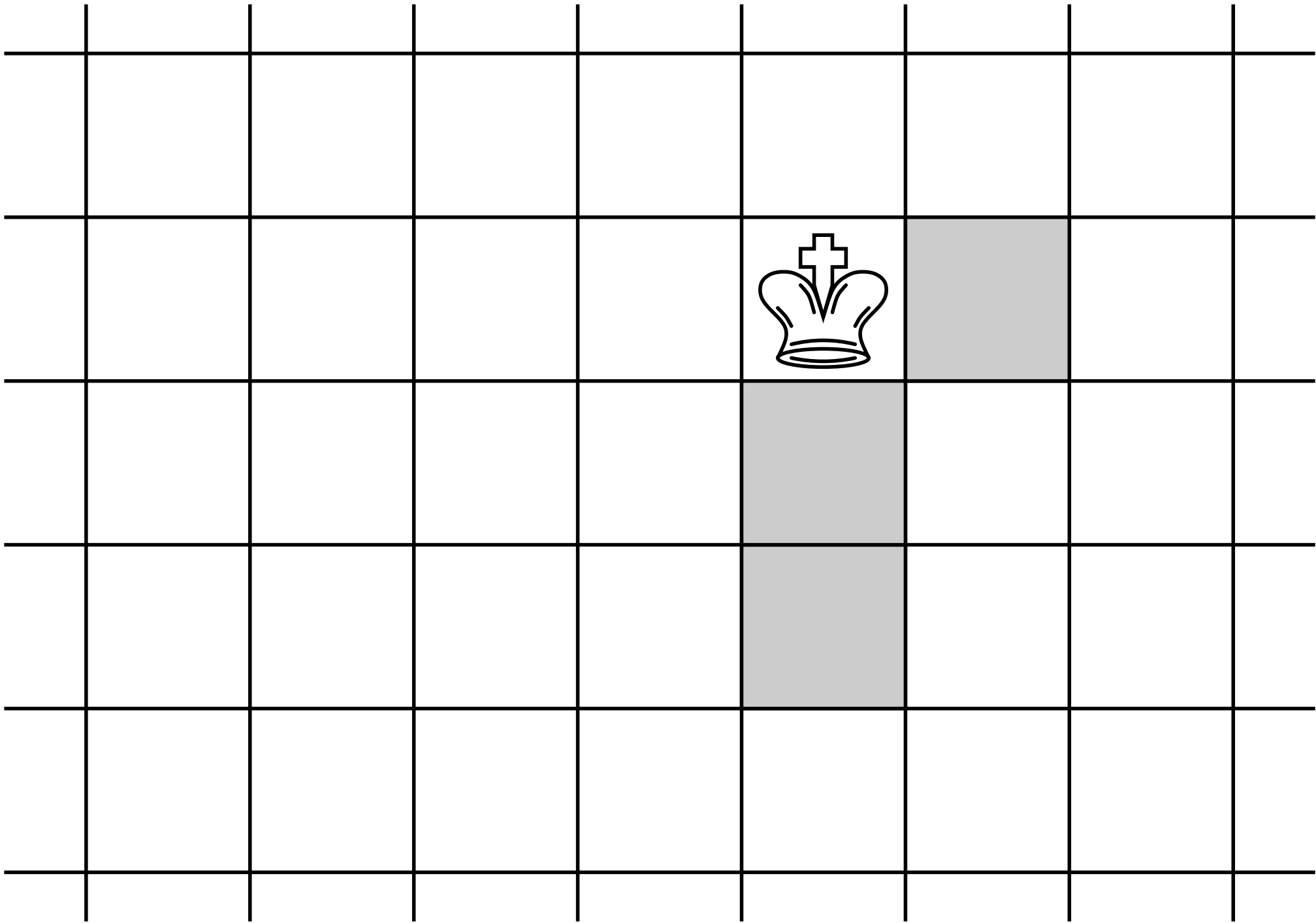


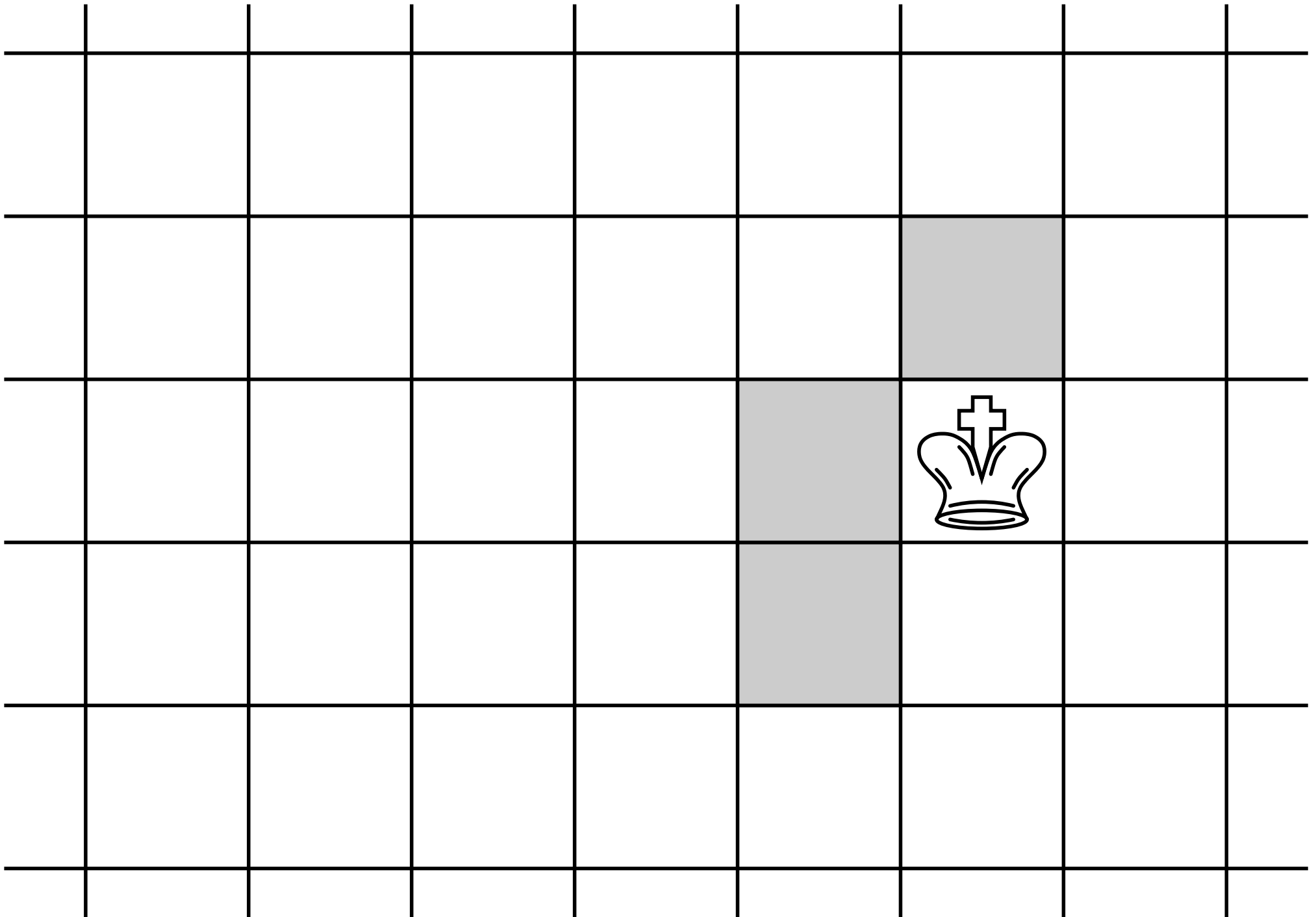


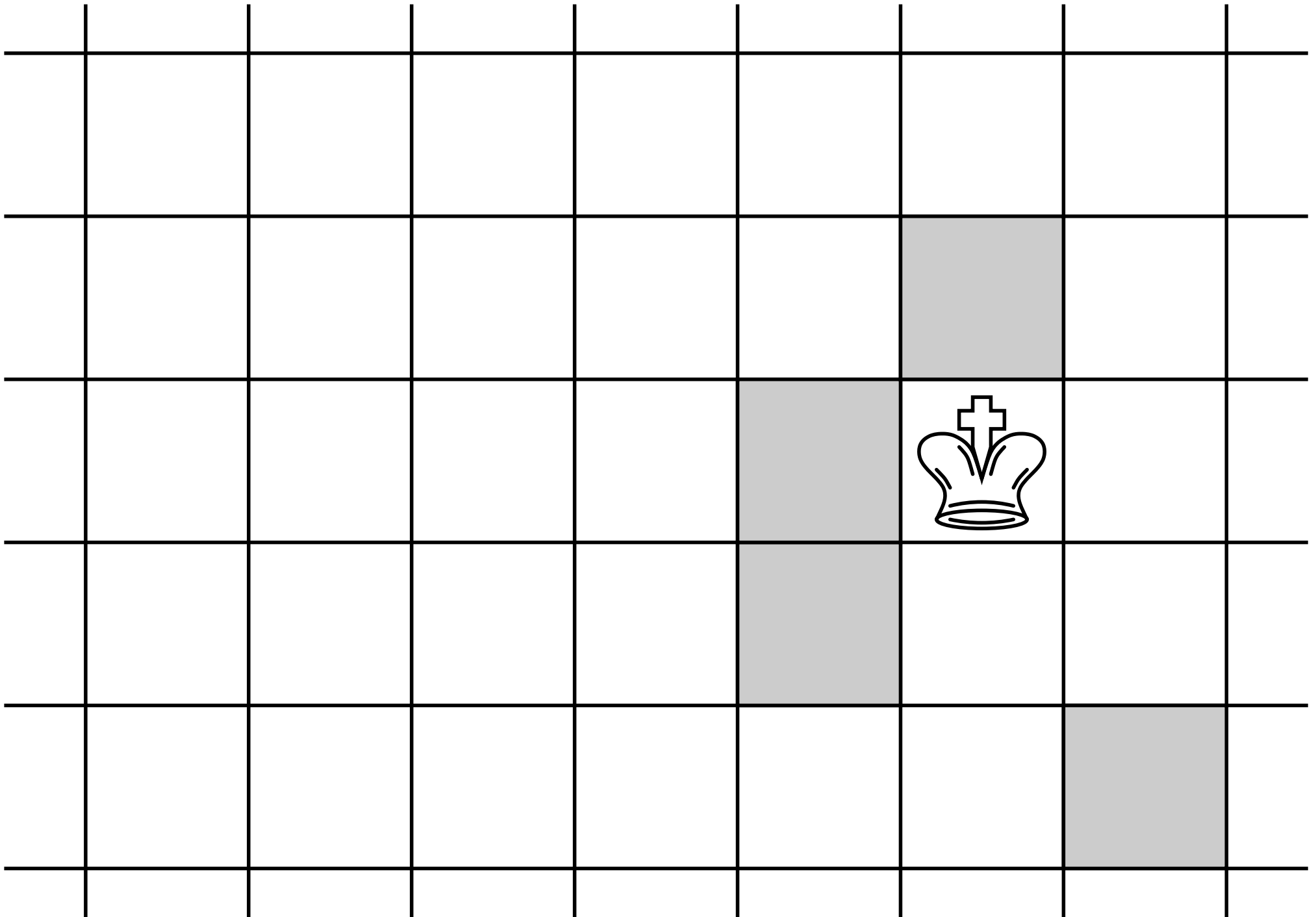


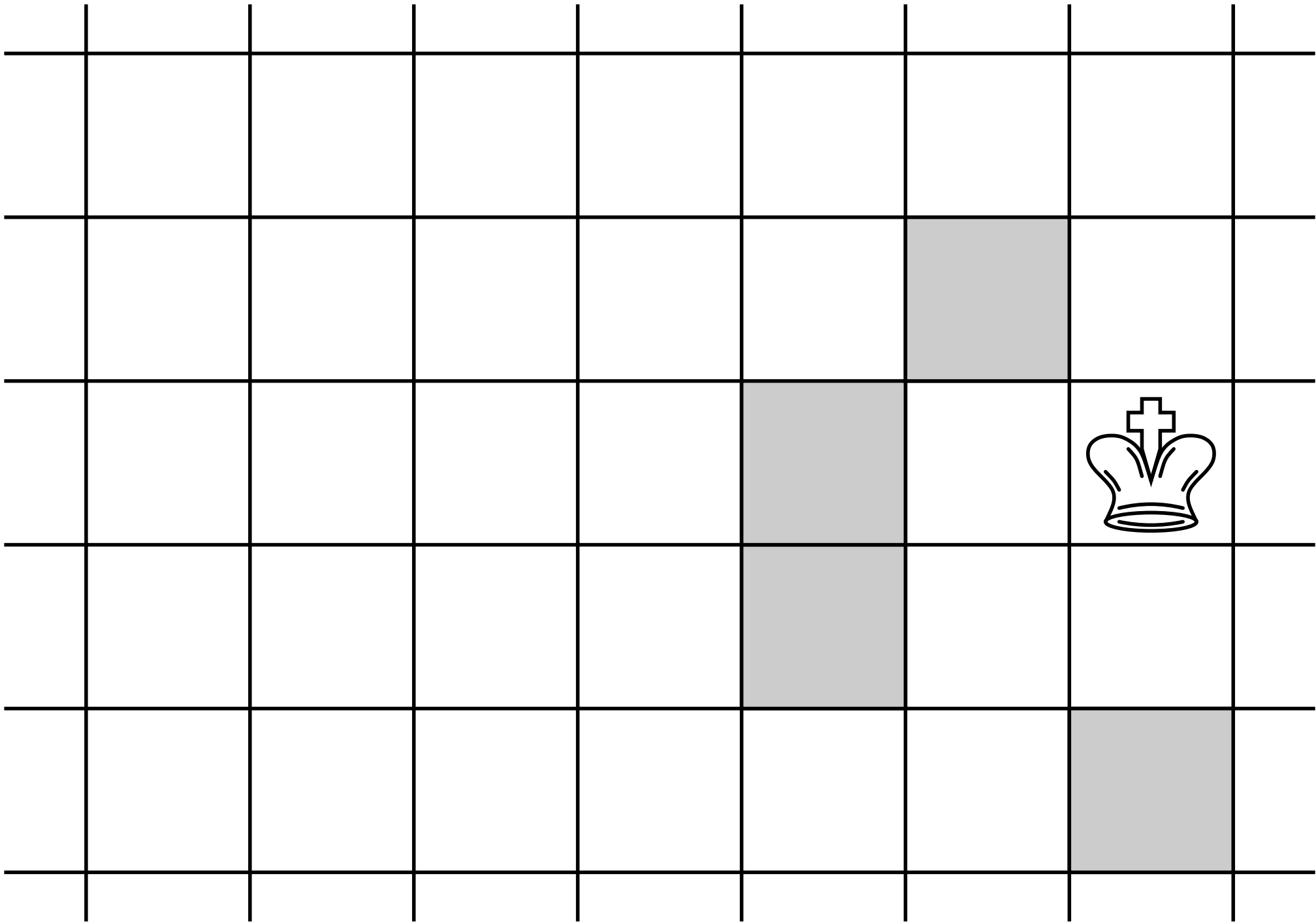


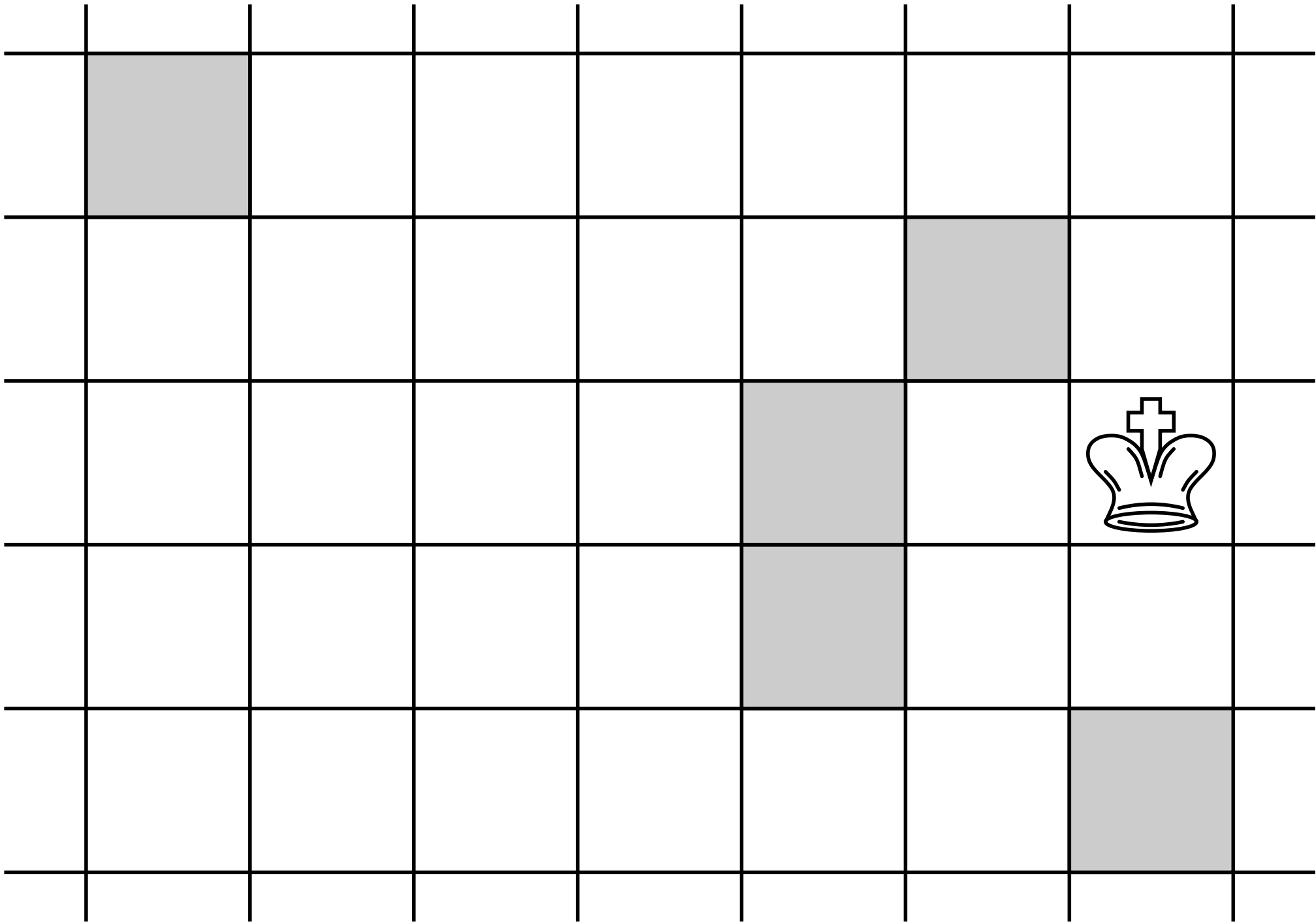










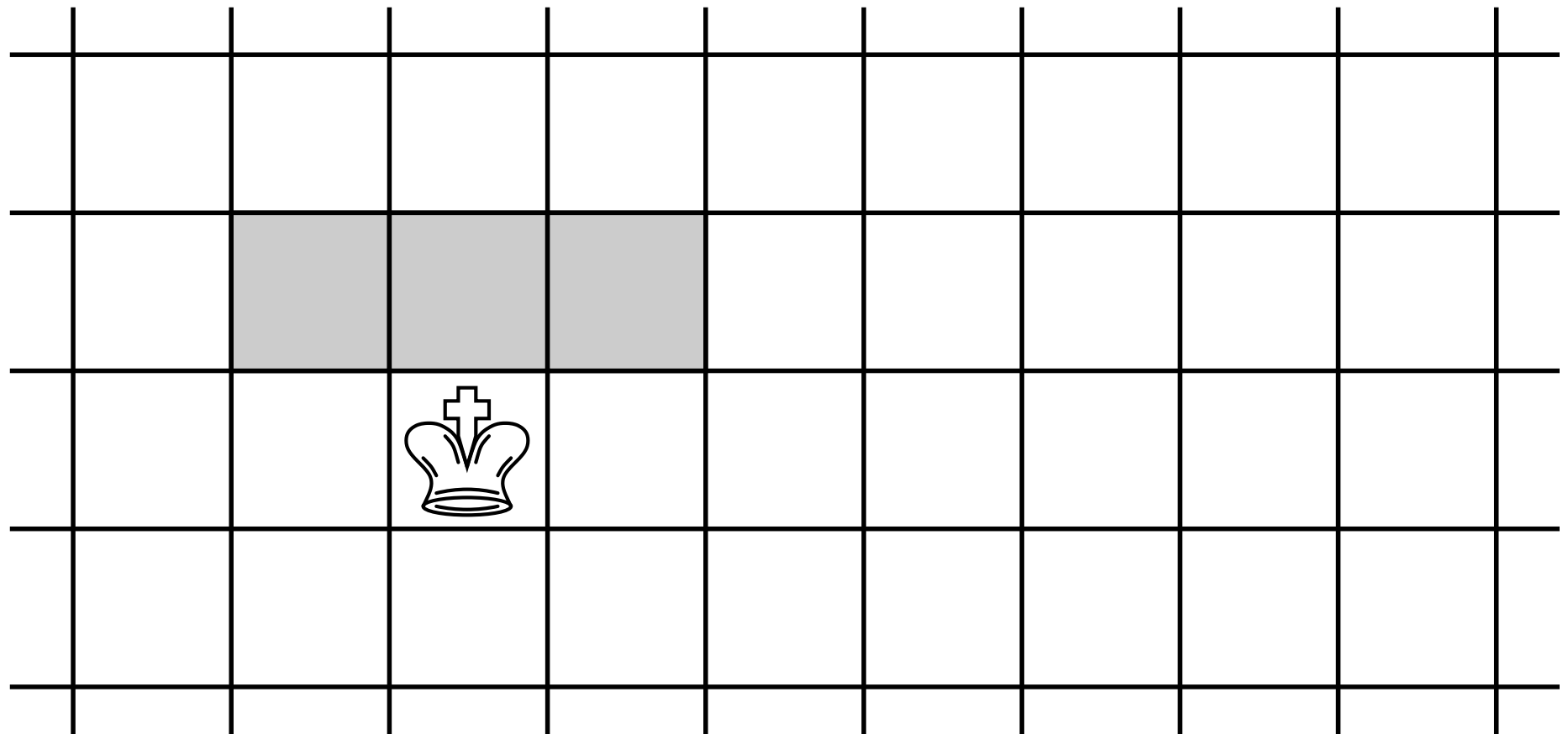


Theorem [Berlekamp]

The chess king can be caught on an infinite checkers board.

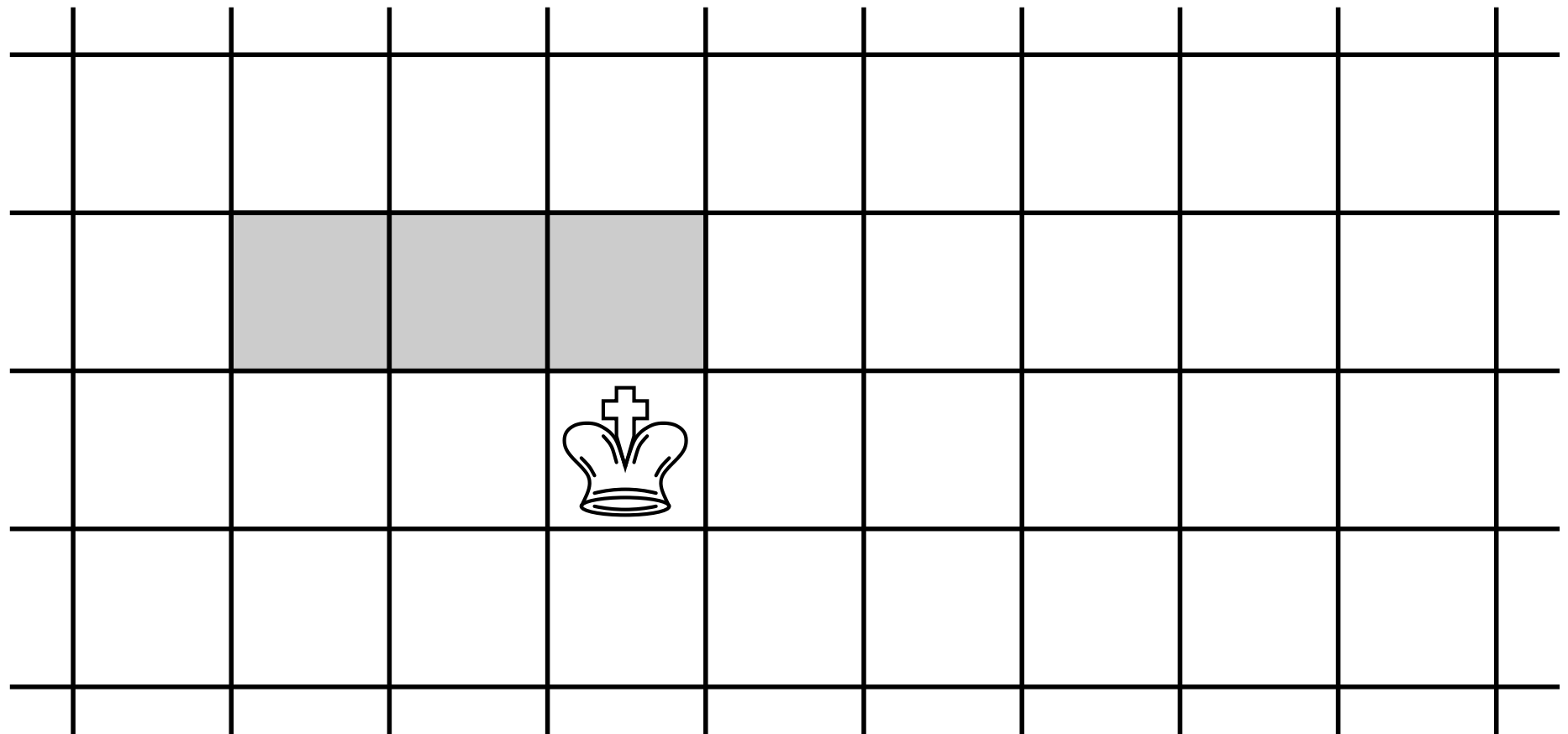
Theorem [Berlekamp]

The chess king can be caught on an infinite checkers board.



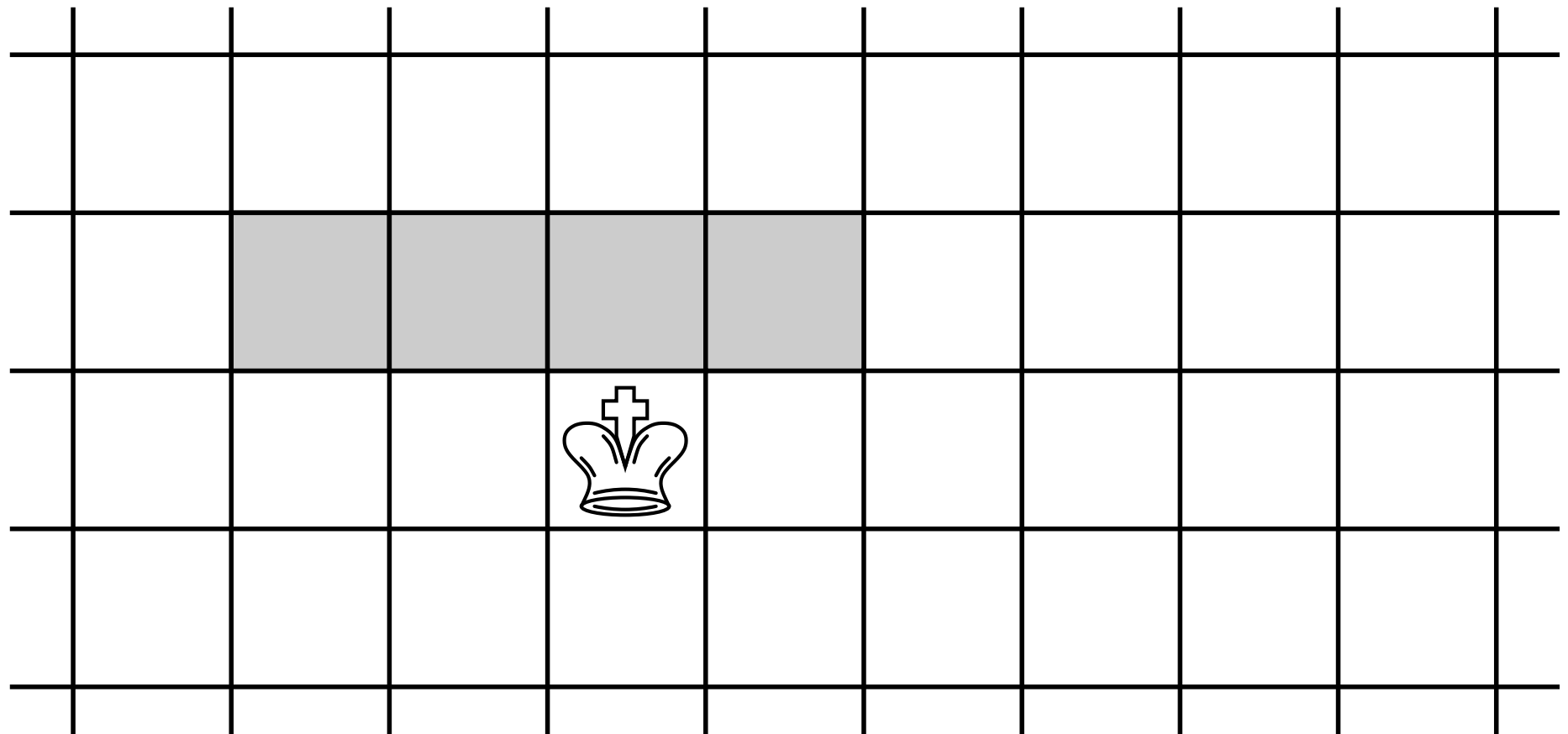
Theorem [Berlekamp]

The chess king can be caught on an infinite checkers board.



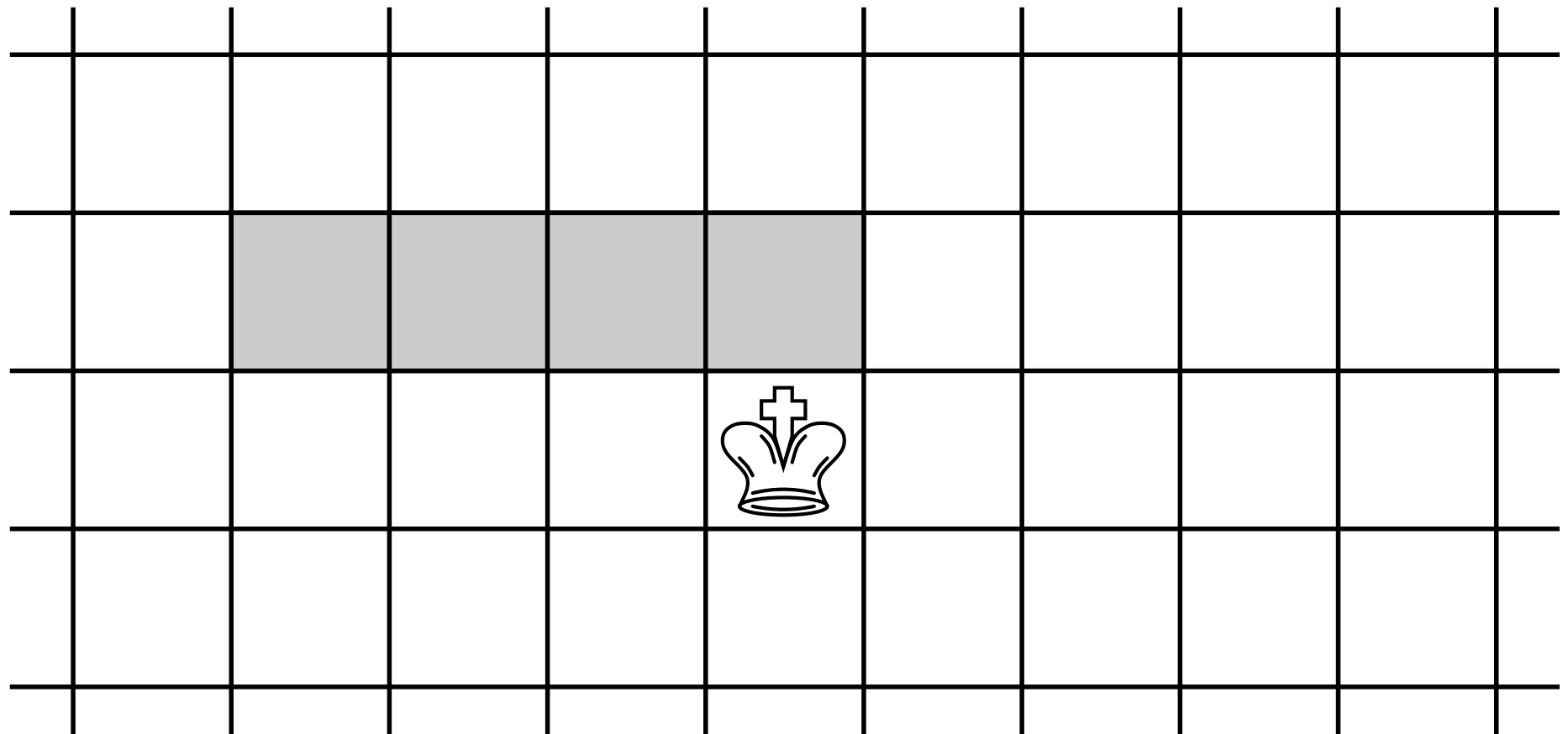
Theorem [Berlekamp]

The chess king can be caught on an infinite checkers board.



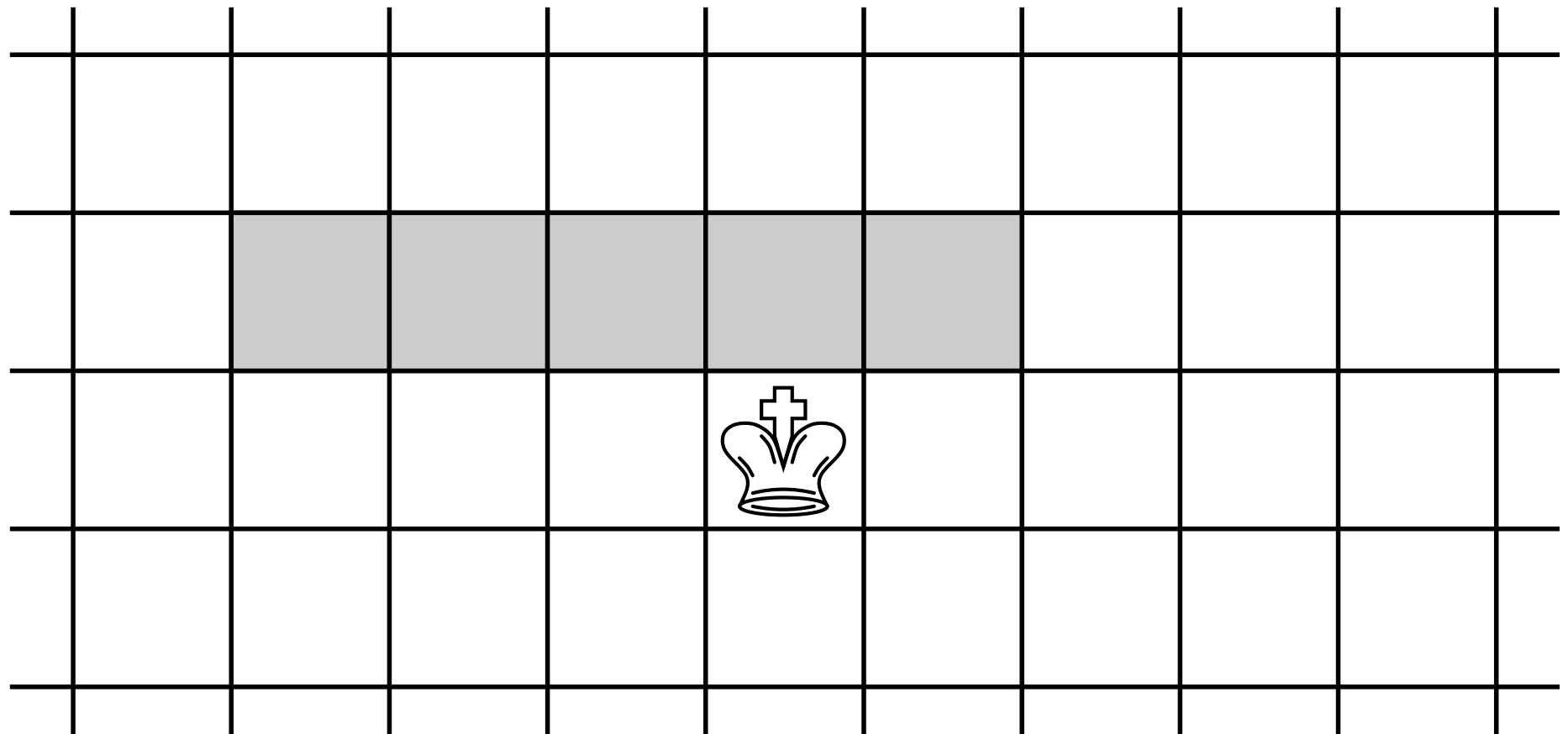
Theorem [Berlekamp]

The chess king can be caught on an infinite checkers board.



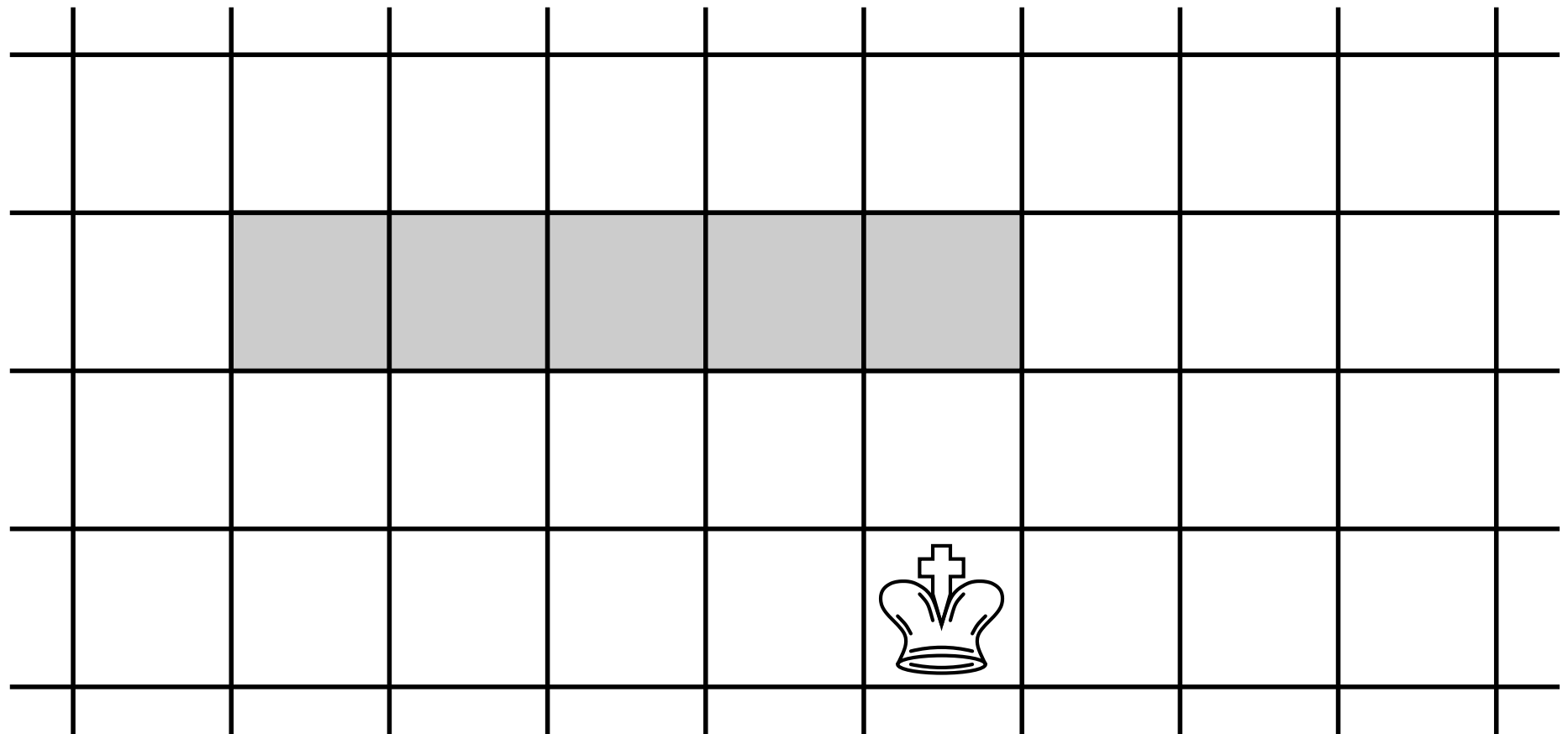
Theorem [Berlekamp]

The chess king can be caught on an infinite checkers board.



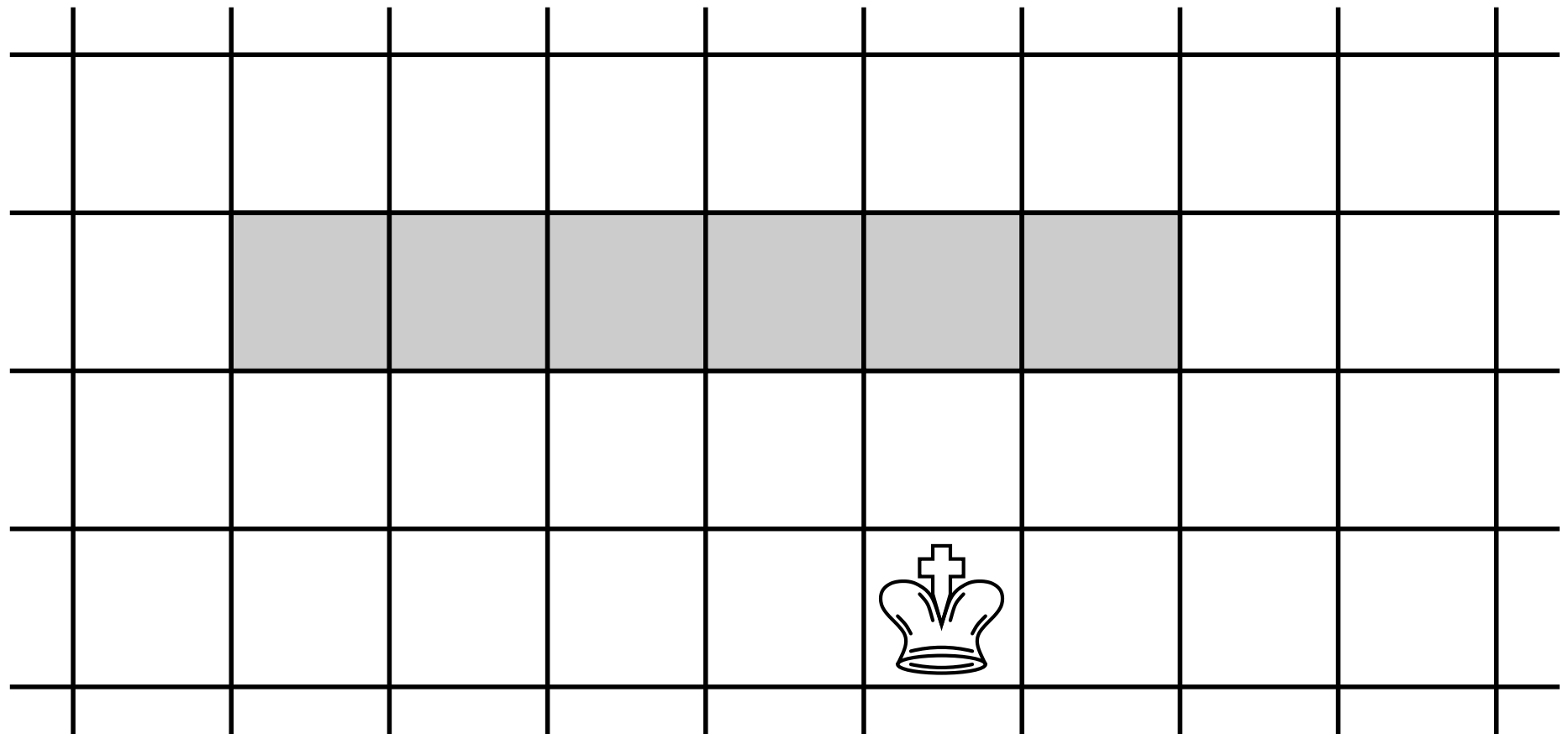
Theorem [Berlekamp]

The chess king can be caught on an infinite checkers board.



Theorem [Berlekamp]

The chess king can be caught on an infinite checkers board.

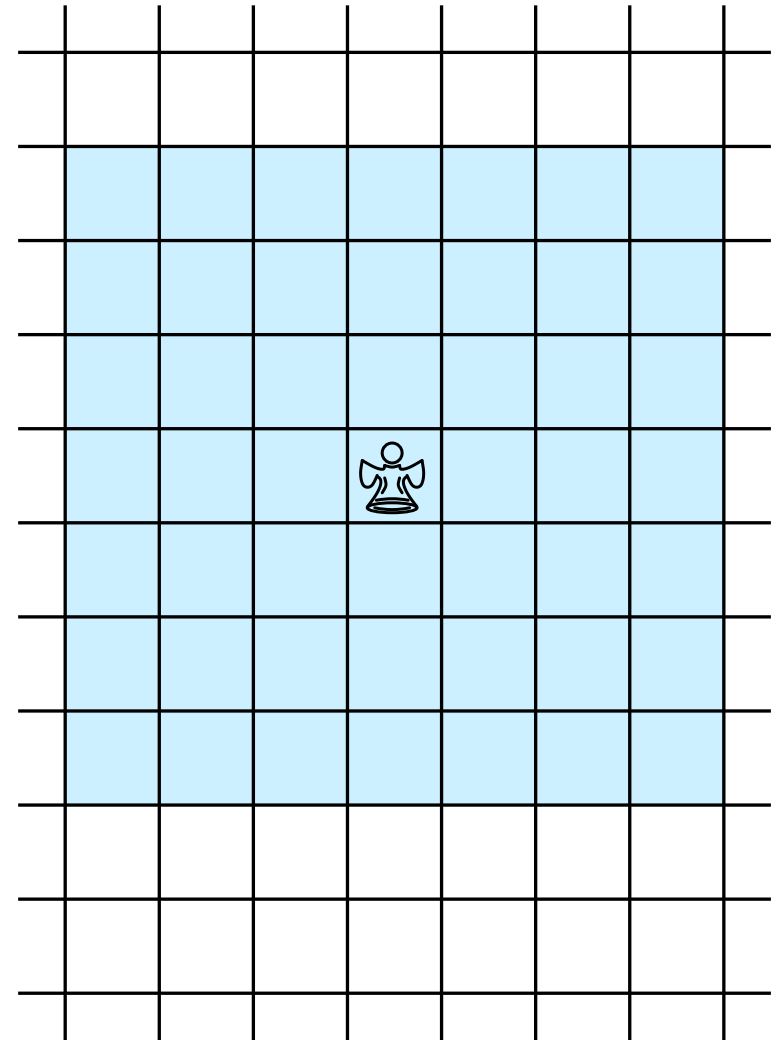


The Angel Problem

Definition

[Berlekamp, Conway, Guy]

A k -Angel can “fly” in one move to any unblocked square at distance at most k .

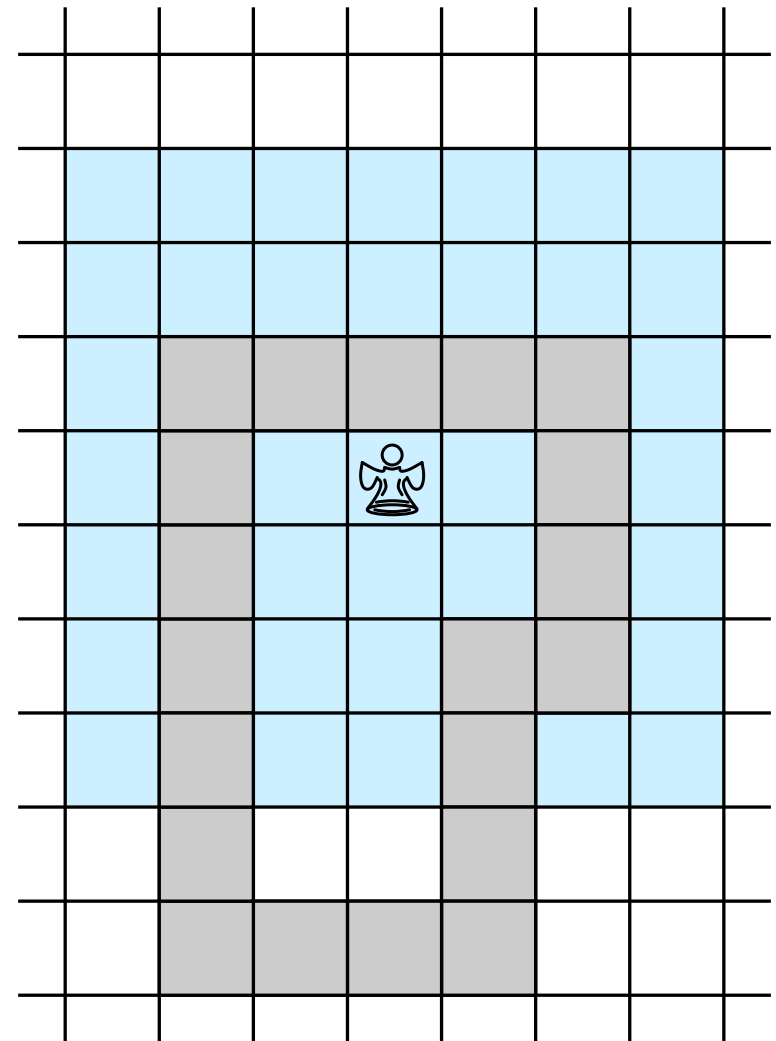


The Angel Problem

Definition

[Berlekamp, Conway, Guy]

A k -Angel can “fly” in one move to any unblocked square at distance at most k .



The Angel Problem

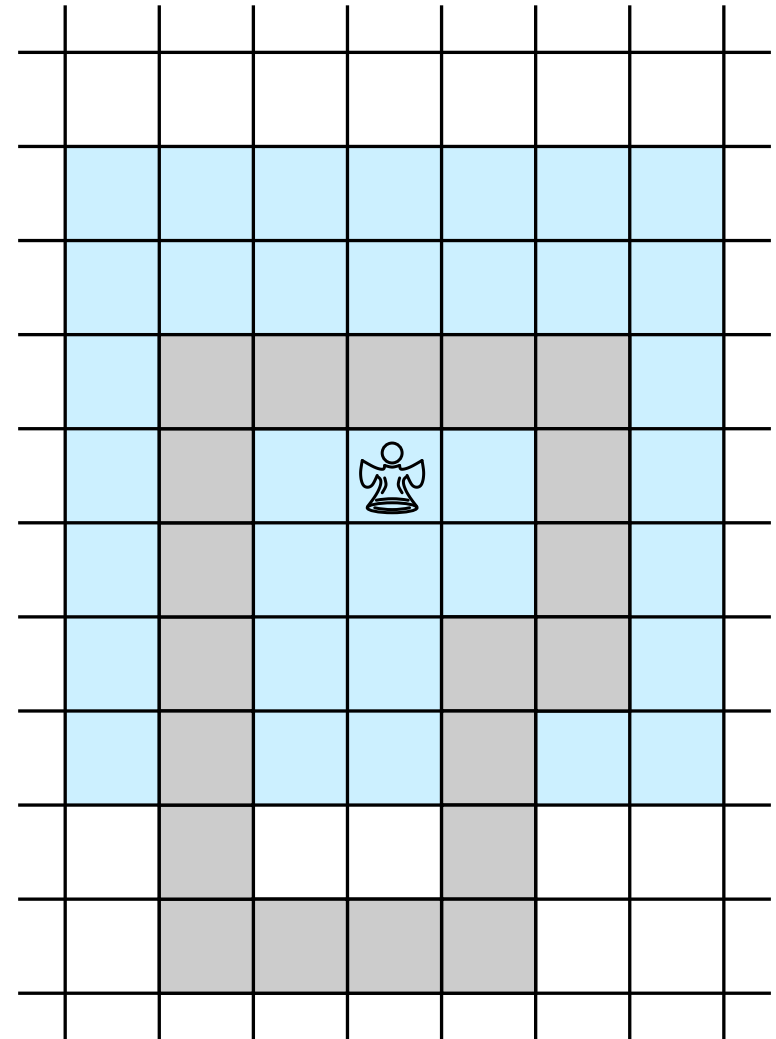
Definition

[Berlekamp, Conway, Guy]

A k -Angel can “fly” in one move to any unblocked square at distance at most k .

Open Problem

Can some k -Angel of some finite power k escape his opponent, the Devil, forever.



Only Fools Rush in

Definition

A **Fool** is an Angel who commits himself to increasing his **y**-coordinate in every move.

Only Fools Rush in

Definition

A **Fool** is an Angel who commits himself to increasing his **y**-coordinate in every move.

Theorem [Conway]

The Devil catches any **k**-Fool of finite power **k**.

Between 1-Angel and 2-Angel

Only the destiny of the 1-Angel (= chess king) is known.

For all other k -Angels, $k \geq 2$, the outcome is open.

Between 1-Angel and 2-Angel

Only the destiny of the 1-Angel (= chess king) is known.

For all other k -Angels, $k \geq 2$, the outcome is open.

We don't even know whether the chess knight can be caught.

Between 1-Angel and 2-Angel

Only the destiny of the 1-Angel (= chess king) is known.

For all other k -Angels, $k \geq 2$, the outcome is open.

We don't even know whether the chess knight can be caught.

Observation: The 2-Angel is actually $4\times$ stronger than the 1-Angel.
(double speed and double-width obstacles)

Between 1-Angel and 2-Angel

Only the destiny of the 1-Angel (= chess king) is known.

For all other k -Angels, $k \geq 2$, the outcome is open.

We don't even know whether the chess knight can be caught.

Observation: The 2-Angel is actually $4\times$ stronger than the 1-Angel.
(double speed and double-width obstacles)

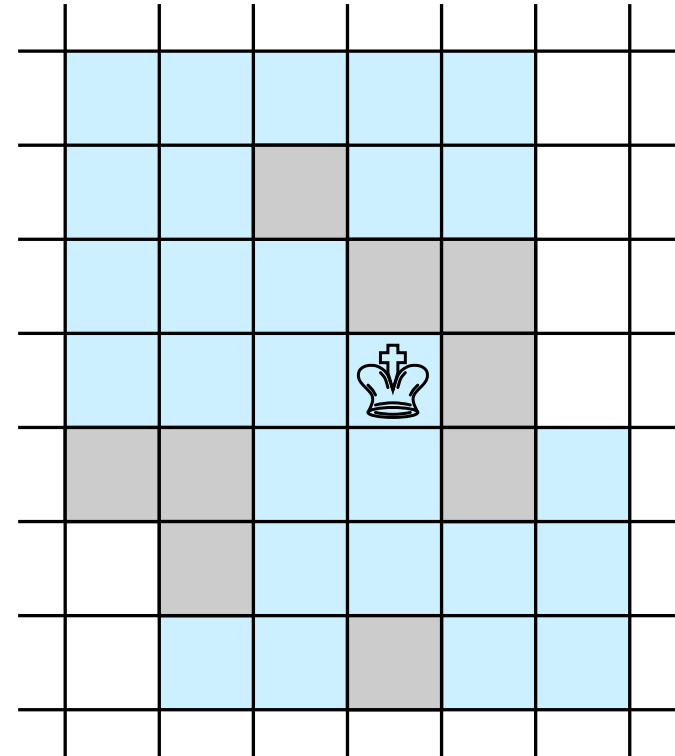
We modify the problem to have **speed** as the only parameter.

Angels With Broken Wings

Deprive Angels of their ability to fly across obstacles.

Definition

A k -King is a k -Angel who can only run, not fly. In each turn he makes k ordinary chess-king moves.



Angels With Broken Wings

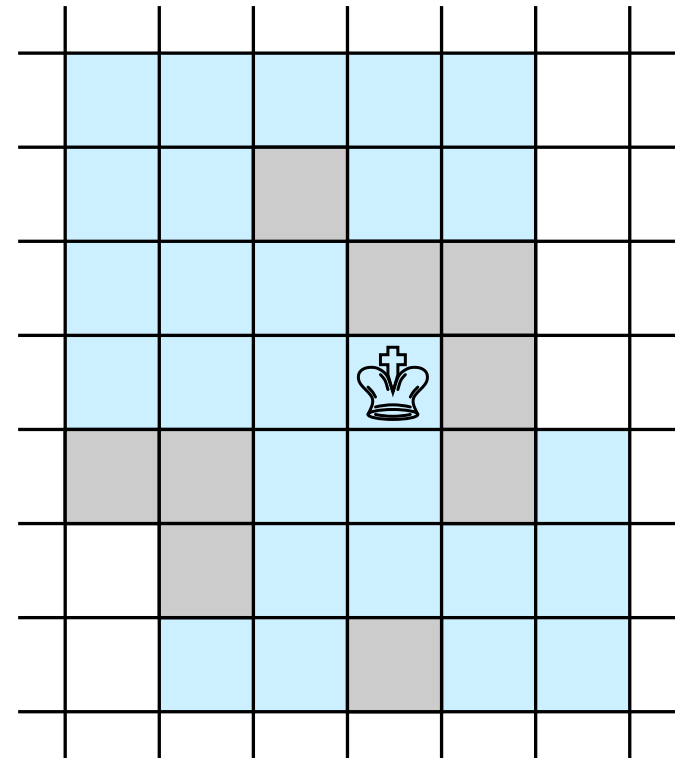
Deprive Angels of their ability to fly across obstacles.

Definition

A k -King is a k -Angel who can only run, not fly. In each turn he makes k ordinary chess-king moves.

Proposition

If the k -Angel can escape forever then so can the $99k^2$ -King.



The Main Result

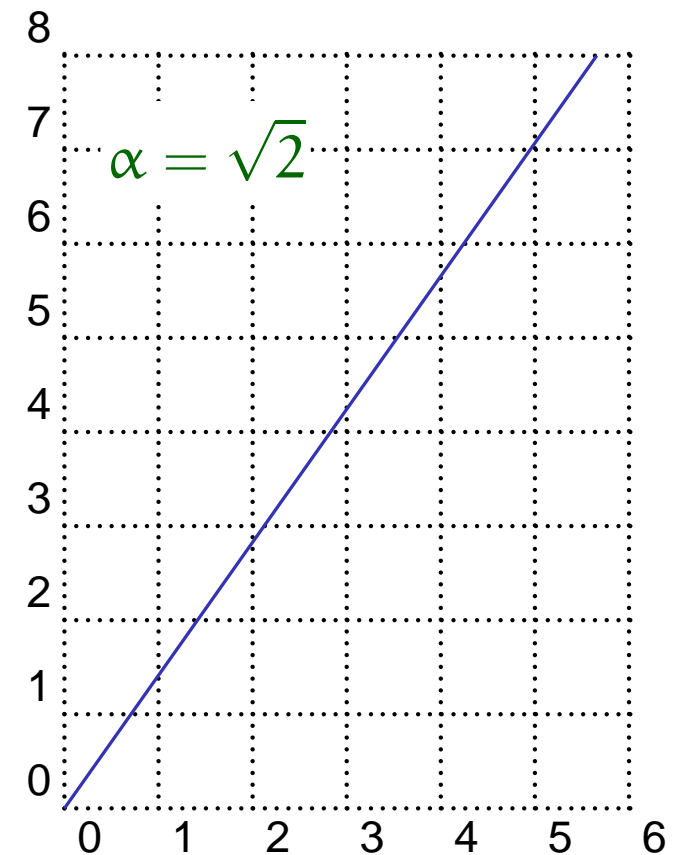
Theorem The Devil can catch any α -King with $\alpha < 2$.

The Main Result

Theorem The Devil can catch any α -King with $\alpha < 2$.

For fractional and irrational speed $\alpha > 1$ define Angel/Devil turns by means of **sturmian sequences**:

Shoot a ray of slope α from the origin and mark crossings with the integer grid:



The Main Result

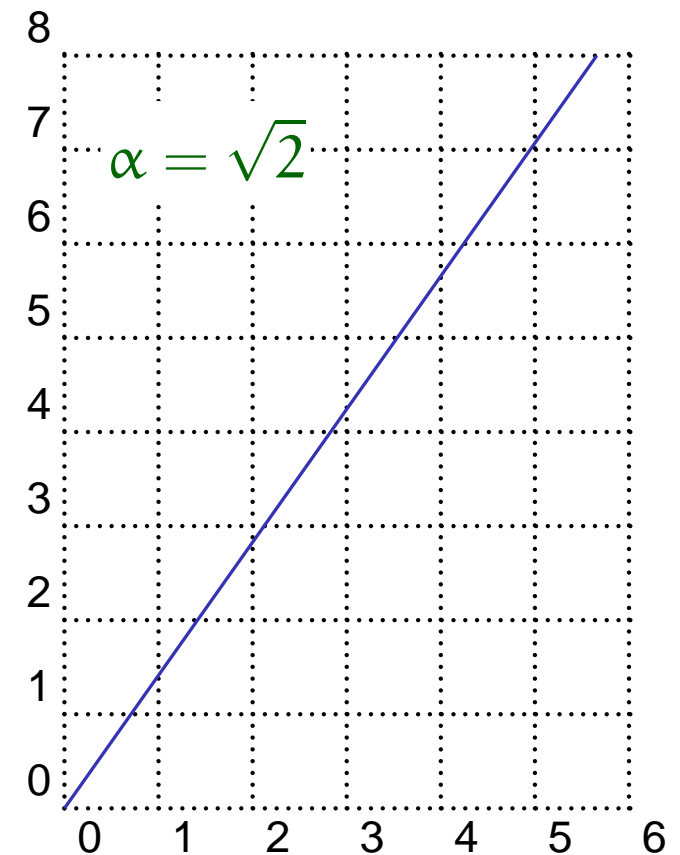
Theorem The Devil can catch any α -King with $\alpha < 2$.

For fractional and irrational speed $\alpha > 1$ define Angel/Devil turns by means of **sturmian sequences**:

Shoot a ray of slope α from the origin and mark crossings with the integer grid:

horizontal line \rightarrow King step

vertical line \rightarrow Devil move



The Main Result

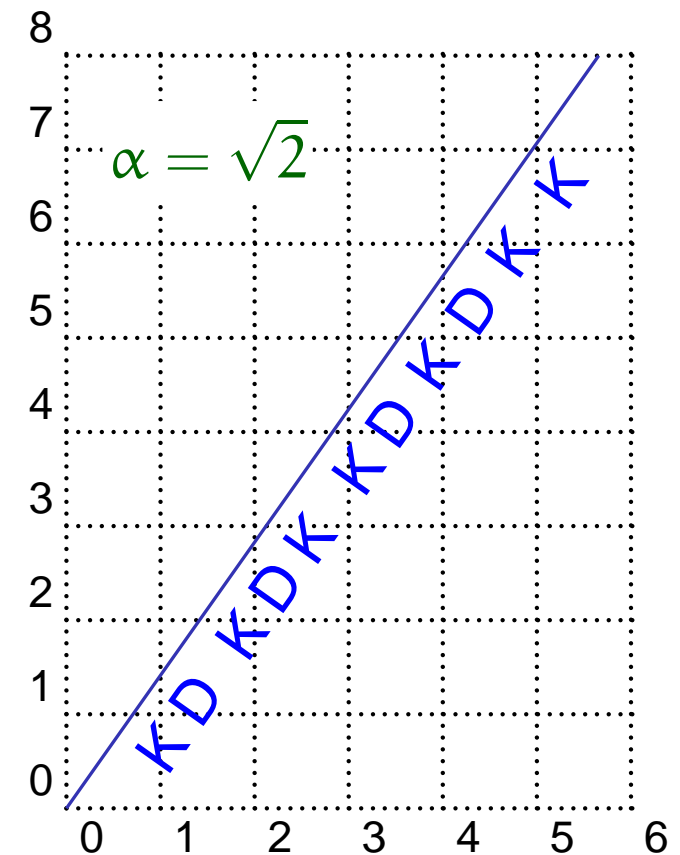
Theorem The Devil can catch any α -King with $\alpha < 2$.

For fractional and irrational speed $\alpha > 1$ define Angel/Devil turns by means of sturmian sequences:

Shoot a ray of slope α from the origin and mark crossings with the integer grid:

horizontal line \rightarrow King step

vertical line \rightarrow Devil move



The Main Result

Theorem The Devil can catch any α -King with $\alpha < 2$.

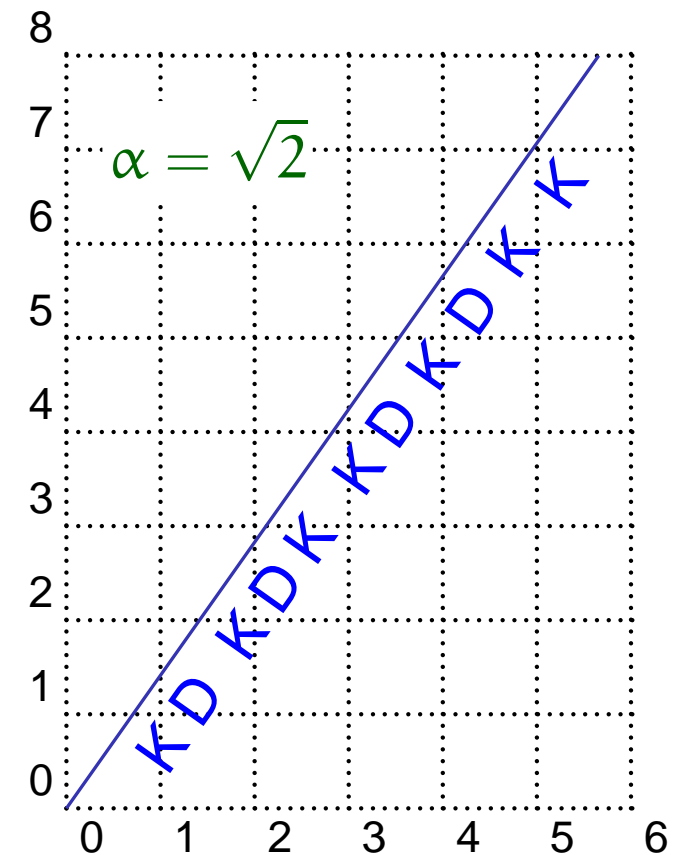
For fractional and irrational speed $\alpha > 1$ define Angel/Devil turns by means of **sturmian sequences**:

Shoot a ray of slope α from the origin and mark crossings with the integer grid:

horizontal line \rightarrow King step

vertical line \rightarrow Devil move

“Lemma.” This distribution is “fair” and shifting of the grid/origin does not affect winning and losing.



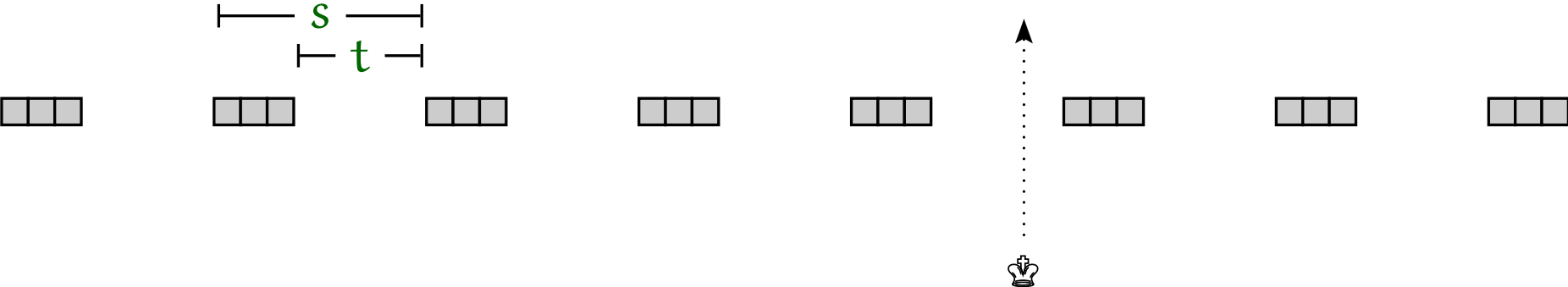
Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly s King moves per t Devil moves.

Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly

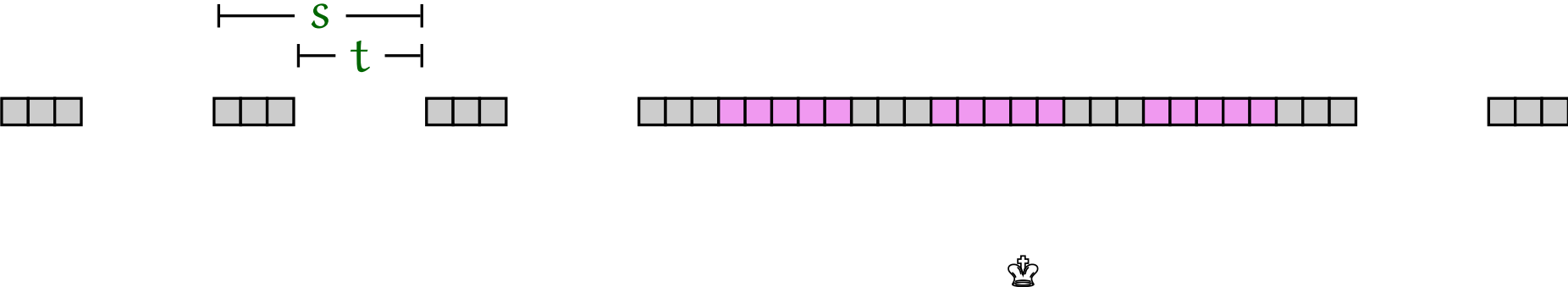
s King moves
per
 t Devil moves.



Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly

s King moves
per
 t Devil moves.



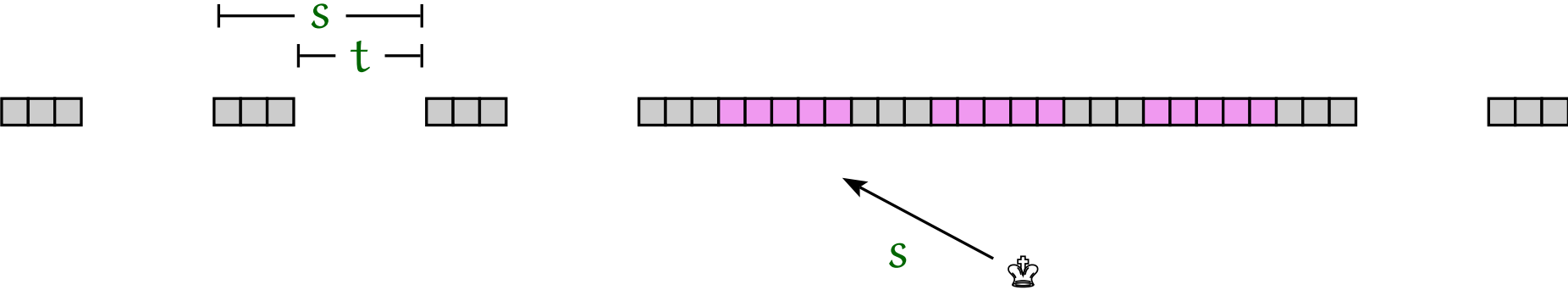
Dynamic Fences

For speed

$$\alpha = \frac{s}{t}$$

we have exactly

s King moves
per
 t Devil moves.



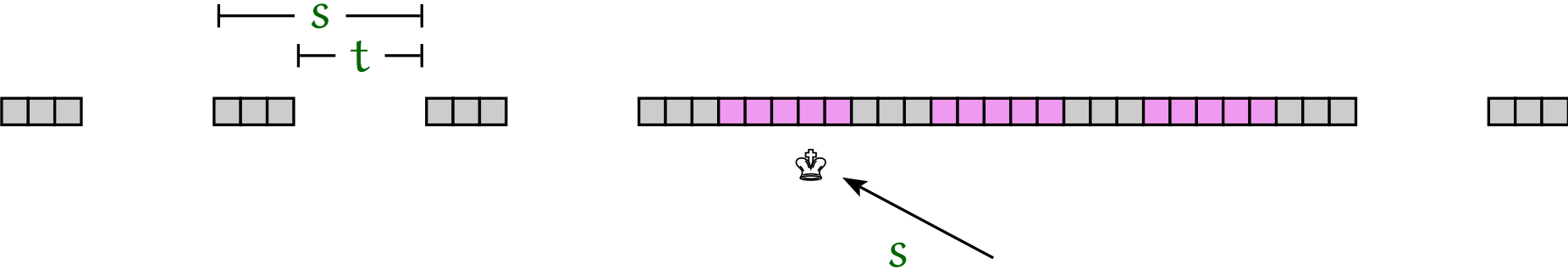
Dynamic Fences

For speed

$$\alpha = \frac{s}{t}$$

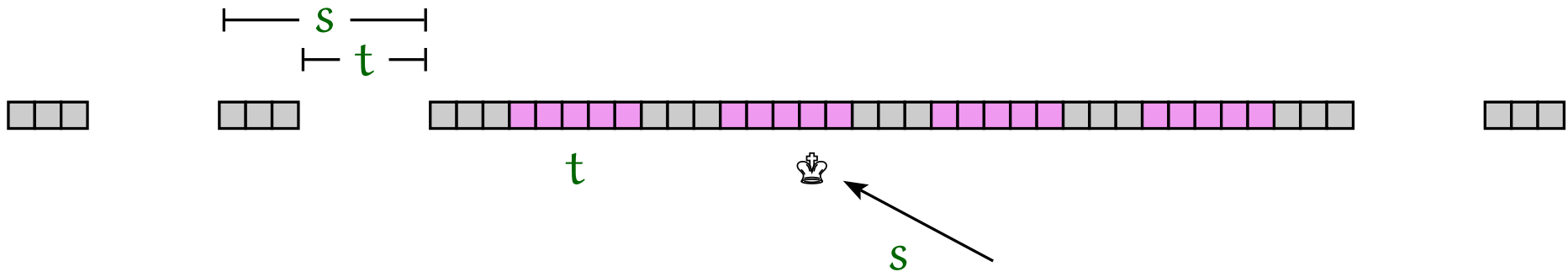
we have exactly

s King moves
per
 t Devil moves.



Dynamic Fences

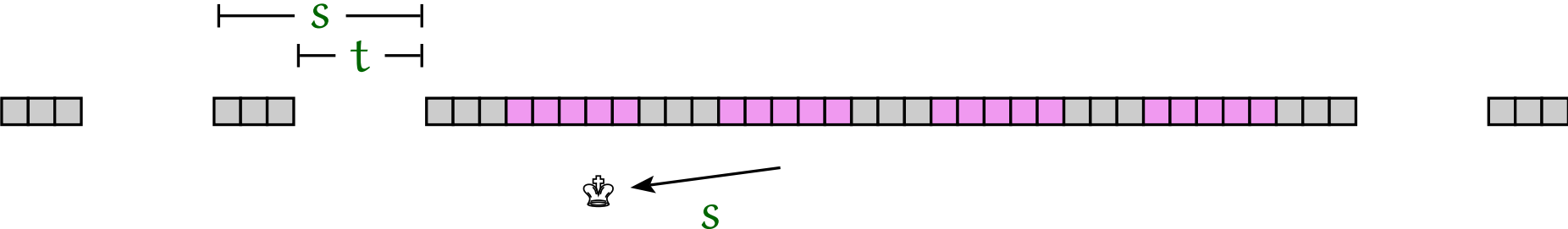
For speed $\alpha = \frac{s}{t}$ we have exactly s King moves per t Devil moves.



Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly

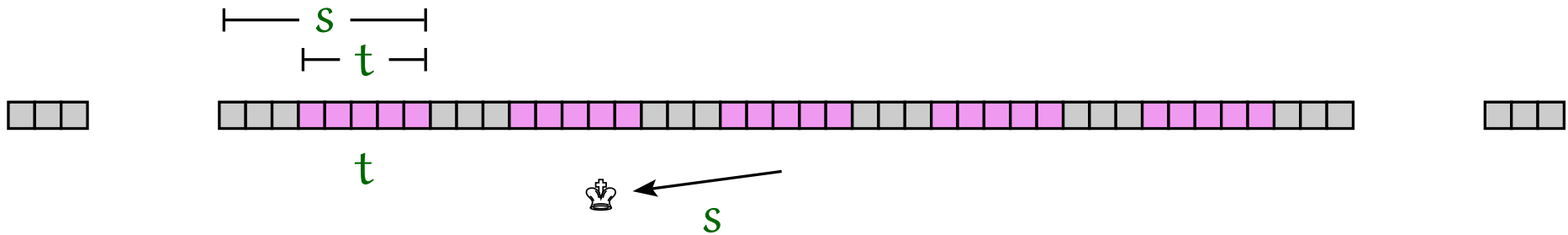
s King moves
per
 t Devil moves.



Dynamic Fences

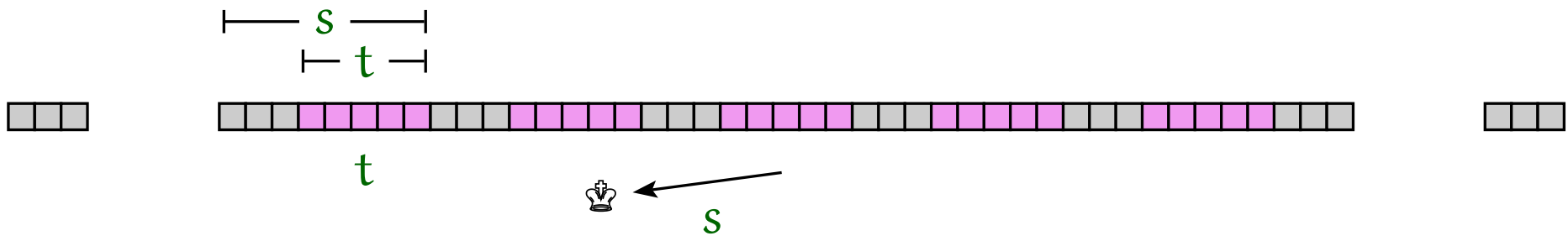
For speed $\alpha = \frac{s}{t}$ we have exactly

s King moves
per
 t Devil moves.



Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly $\frac{s}{t}$ King moves per Devil moves.

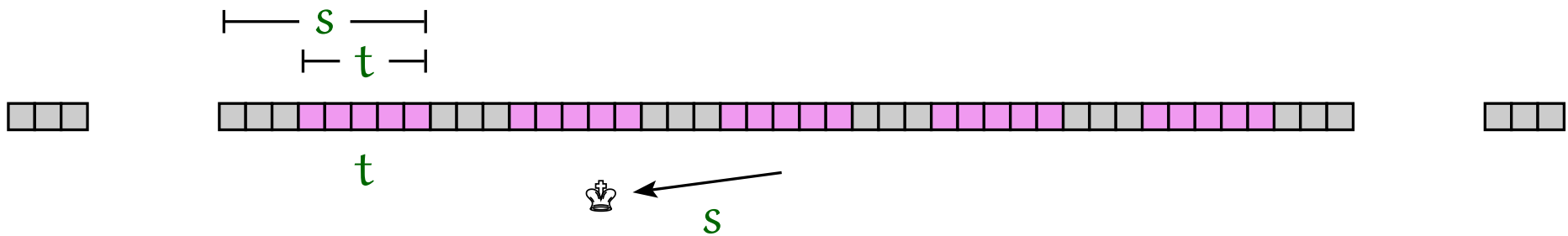


Lemma Against an $\frac{s}{t}$ -King there exist **dynamic fences** of density

$$\frac{s-t}{s} = 1 - \frac{t}{s}$$

Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly s King moves per t Devil moves.

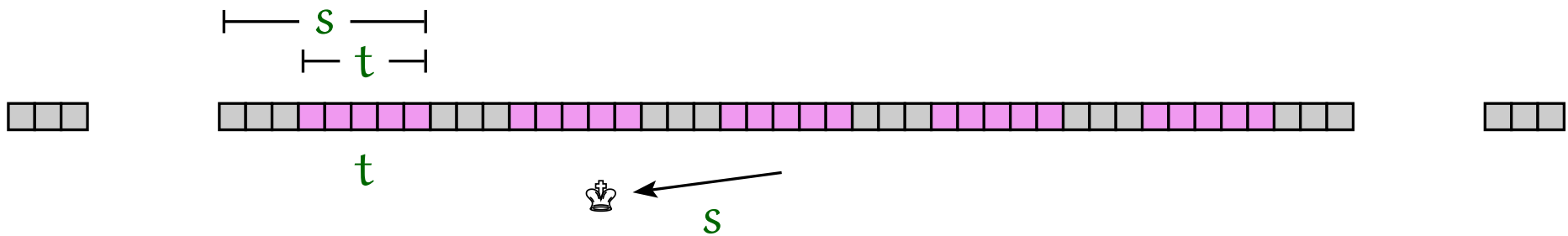


Lemma Against an $\frac{s}{t}$ -King there exist **dynamic fences** of density

$$\frac{s-t}{s} = 1 - \frac{t}{s} < \frac{1}{2} \quad \left(\text{for } \frac{s}{t} < 2 \right)$$

Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly $\frac{s}{t}$ King moves per Devil moves.

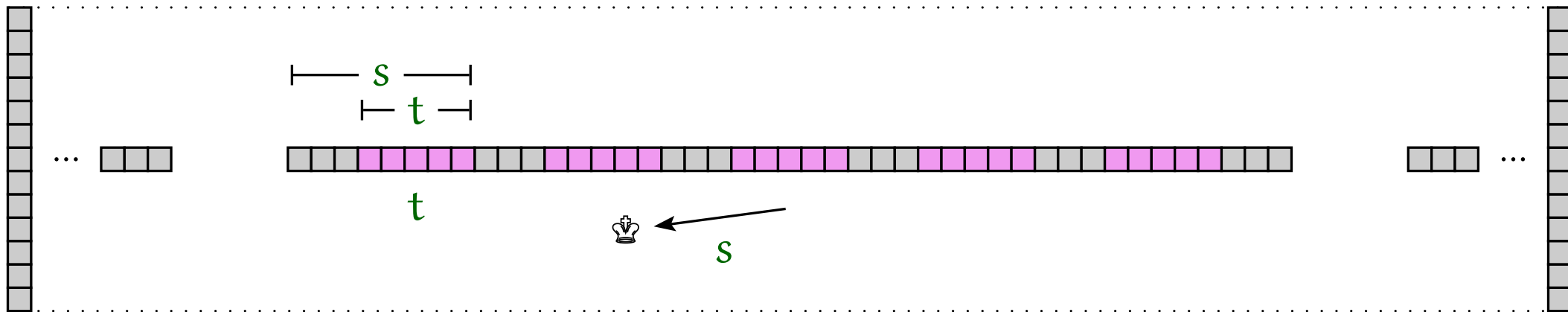


Lemma Against an $\frac{s}{t}$ -King there exist **dynamic fences** of density

$$\frac{s-t}{s} = 1 - \frac{t}{s} < \frac{1}{2} \quad \left(\text{for } \frac{s}{t} < 2 \right)$$

Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly s King moves per t Devil moves.

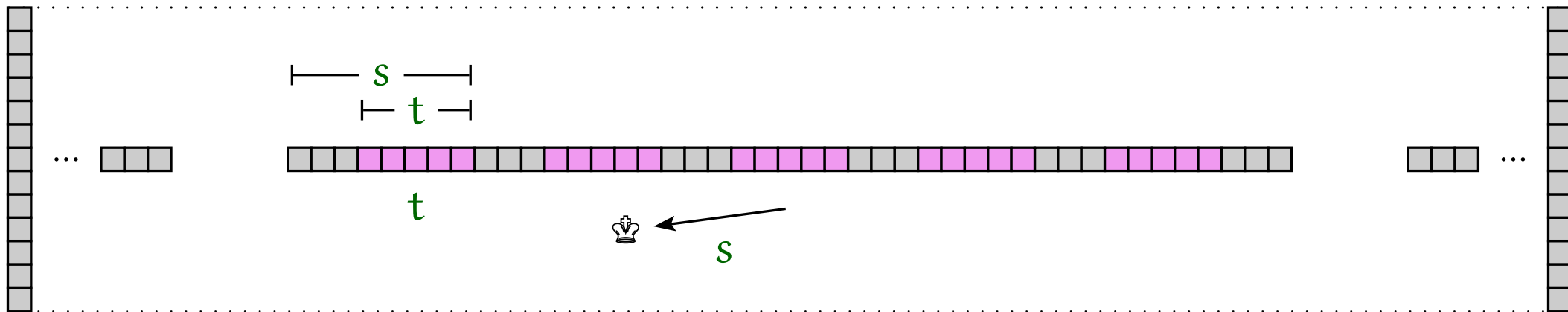


Lemma Against an $\frac{s}{t}$ -King there exist **dynamic fences** of density

$$\frac{s-t}{s} = 1 - \frac{t}{s} < \frac{1}{2} \quad \left(\text{for } \frac{s}{t} < 2 \right)$$

Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly s King moves per t Devil moves.

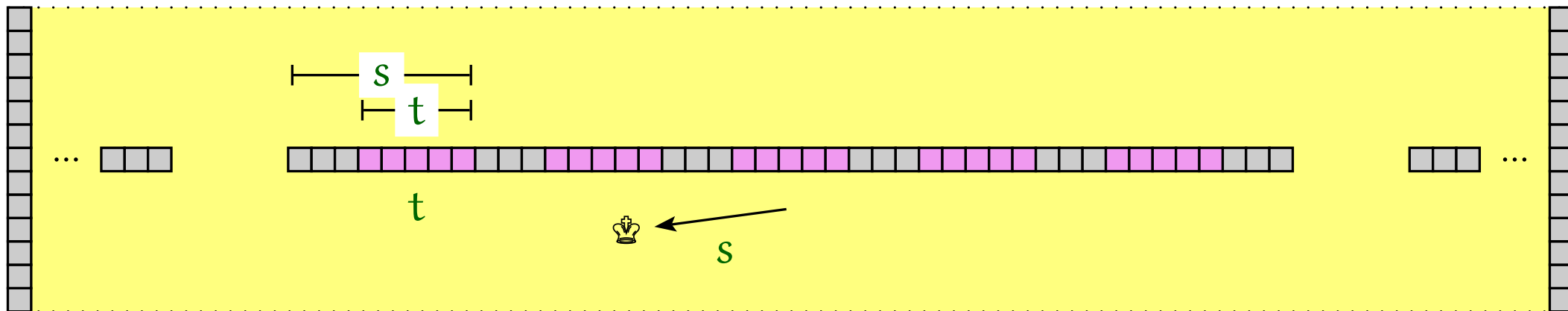


Lemma Against an $\frac{s}{t}$ -King there exist **dynamic fences** of density

$$\frac{s-t}{s} + \varepsilon = 1 - \frac{t}{s} + \varepsilon < \frac{1}{2} \quad \left(\text{for } \frac{s}{t} < 2 \right)$$

Dynamic Fences

For speed $\alpha = \frac{s}{t}$ we have exactly $\frac{s}{t}$ King moves per Devil moves.

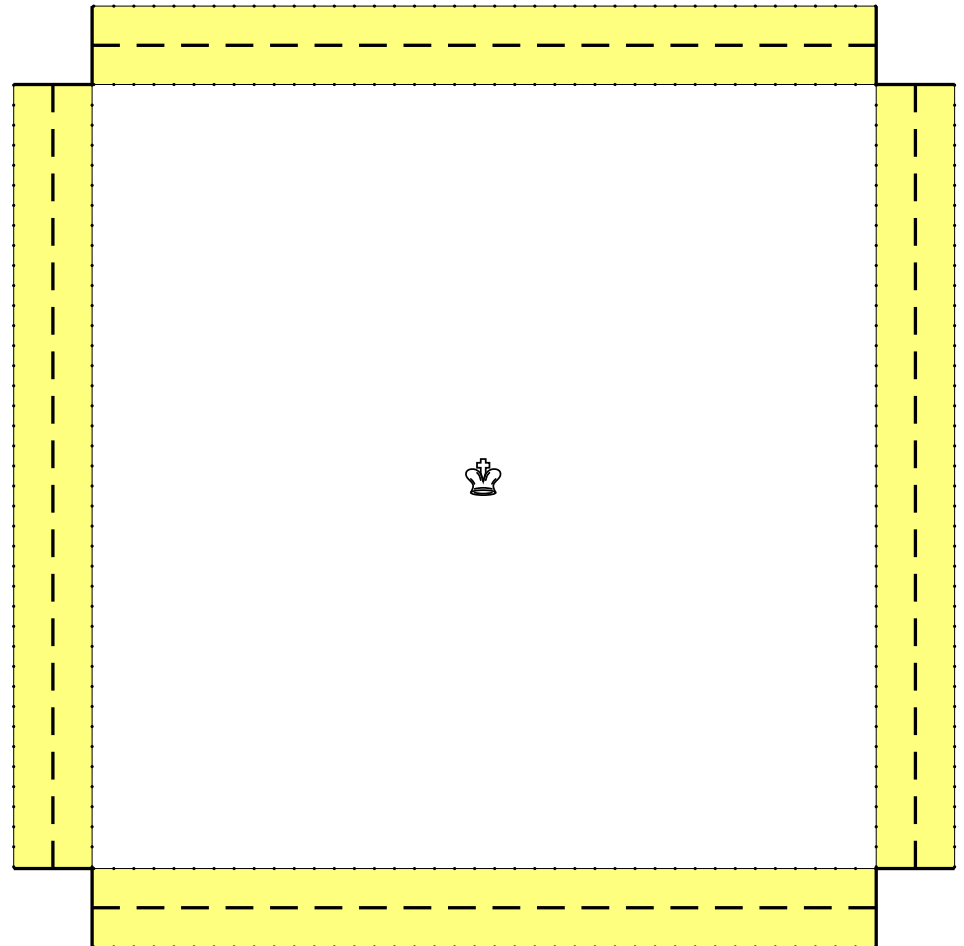


Lemma Against an $\frac{s}{t}$ -King there exist **dynamic fences** of density

$$\frac{s-t}{s} + \varepsilon = 1 - \frac{t}{s} + \varepsilon < \frac{1}{2} \quad \left(\text{for } \frac{s}{t} < 2 \right)$$

Building a Box

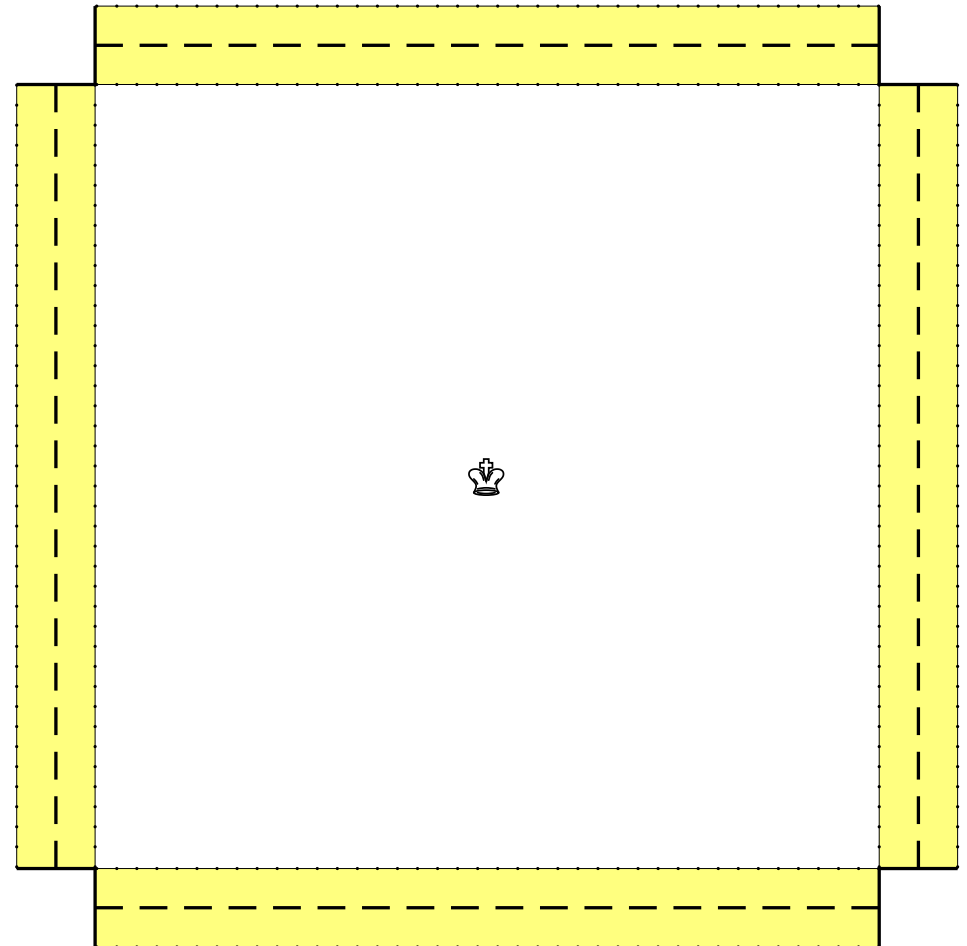
Encircle the King with a box of fences before he can reach the boundary.



Building a Box

Encircle the King with a box of fences before he can reach the boundary.

This only works with fences of *very low* density.



Building a Box

Encircle the King with a box of fences before he can reach the boundary.

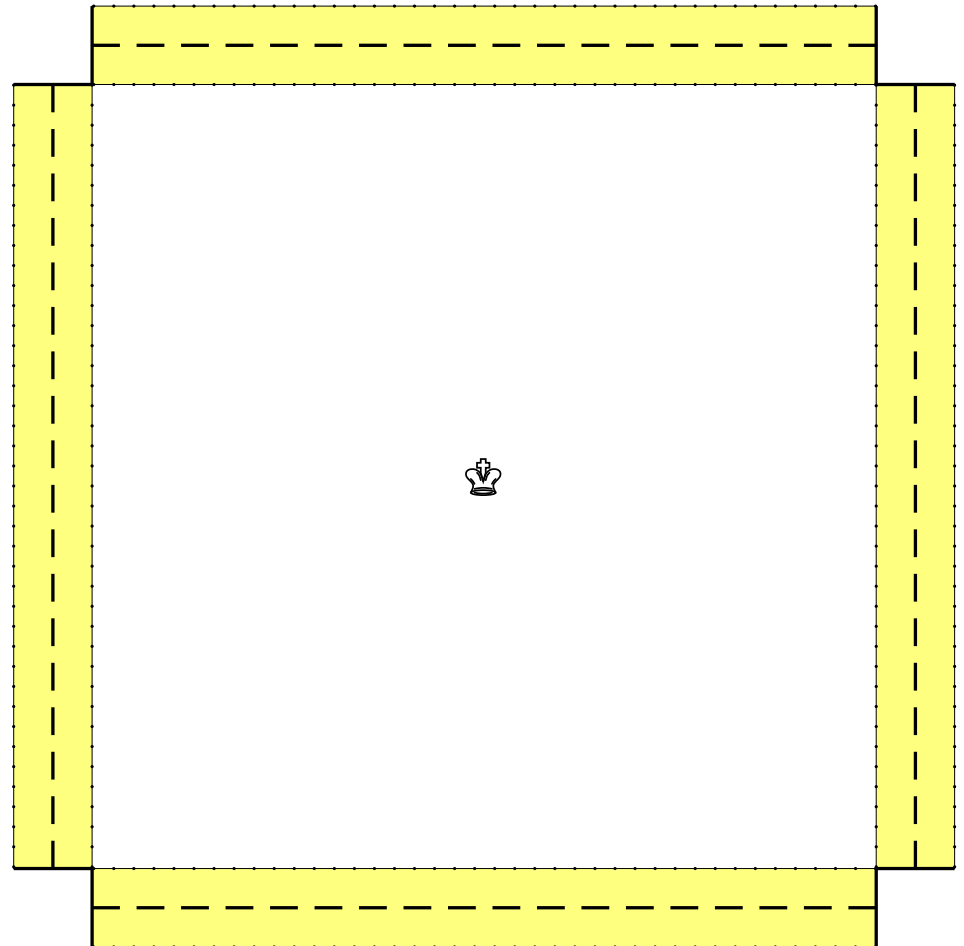
This only works with fences of *very low* density.

A first result:

For $\alpha < 9/8$ we get fences of density $< 1/9$, which the Devil builds

$$\frac{9}{9/8} = 8$$

times faster than the King runs.



Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

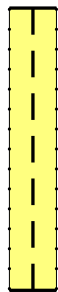
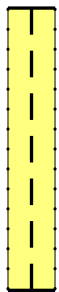
Need smaller densities.

Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences

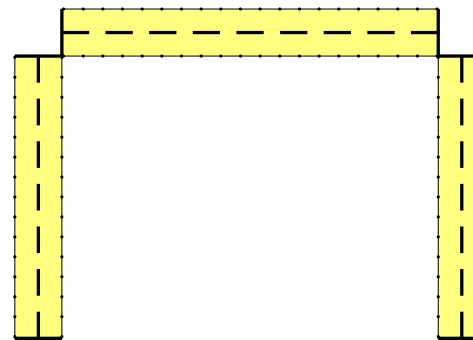
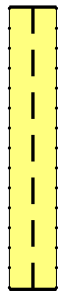
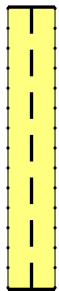


Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences

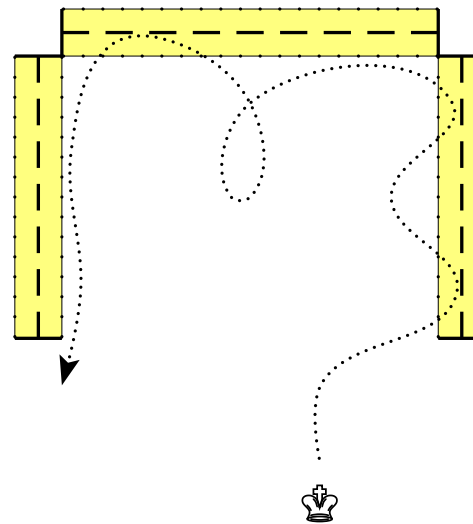
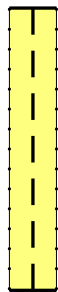
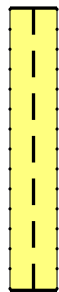


Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences

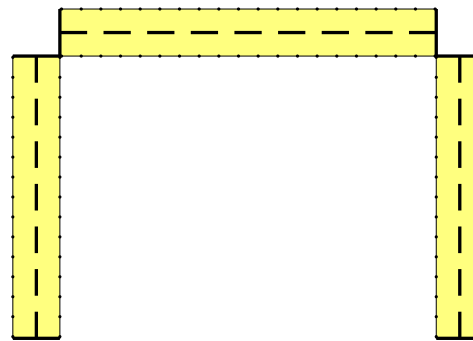
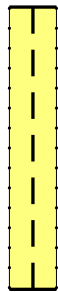
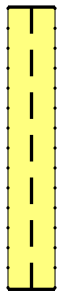


Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences

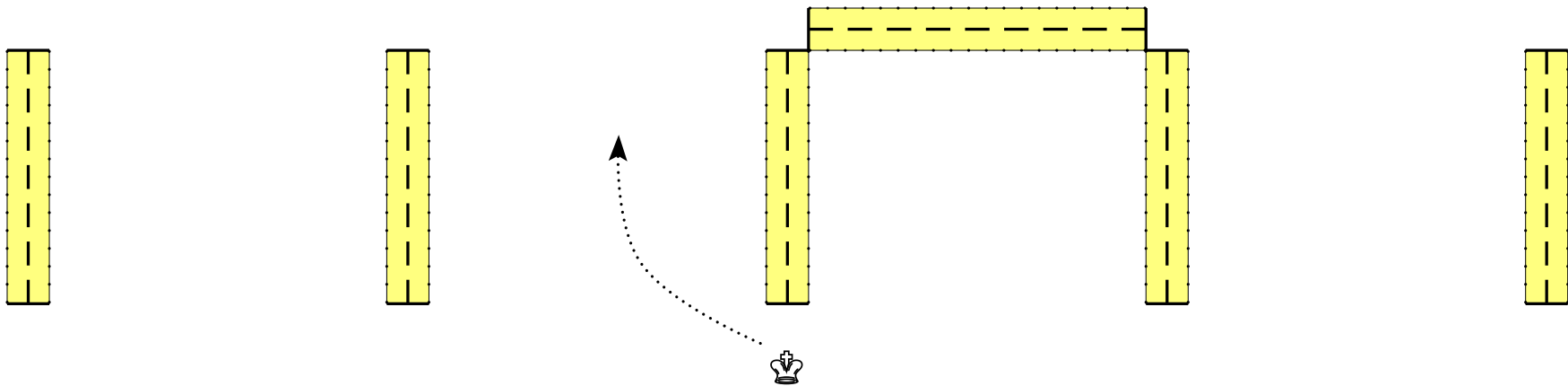


Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences

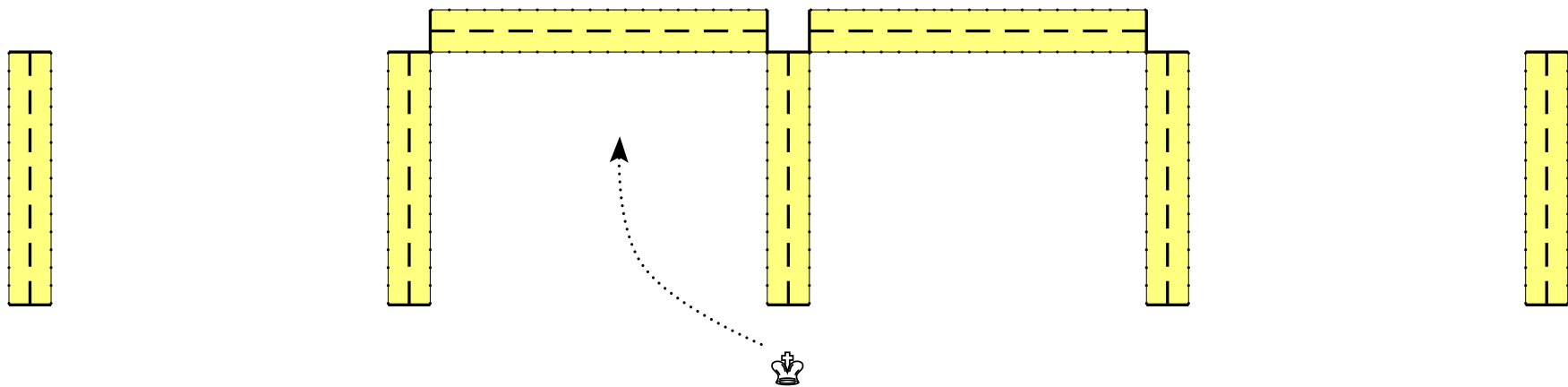


Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences

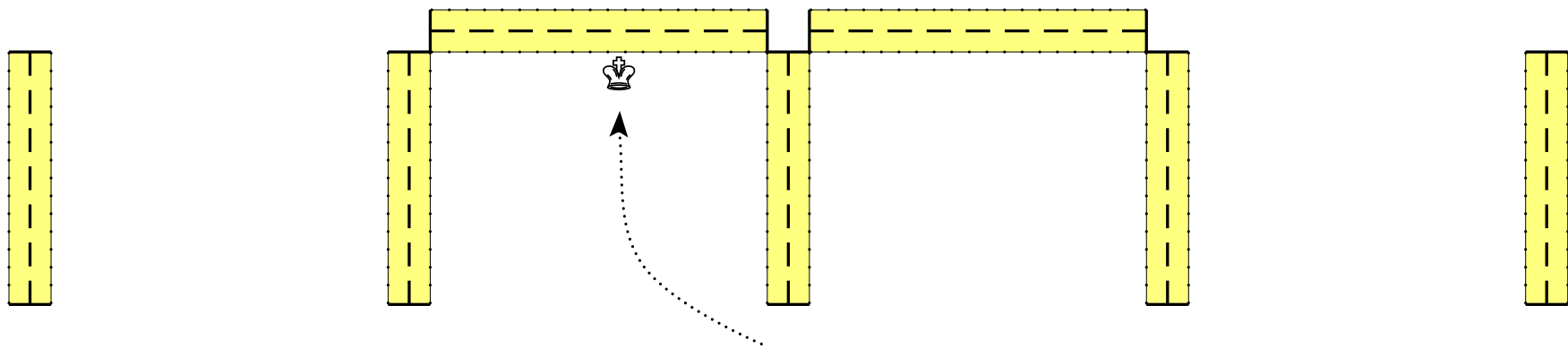


Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences

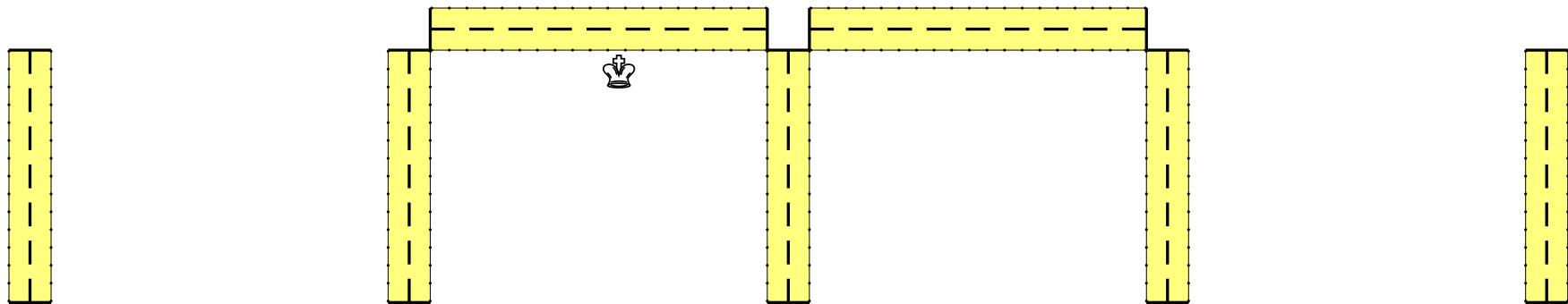


Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences



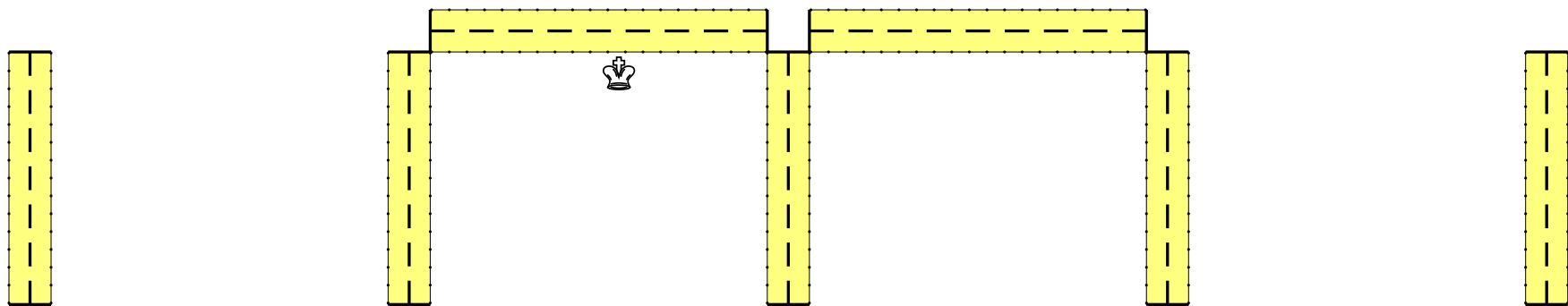
The slots are wider than they are deep, so the total density lies below that of the small fences. (works only for density $< 1/2$)

Smaller Densities

Against King speed $2 - \varepsilon$, fence density $\frac{1}{2} - \varepsilon'$ is not enough.

Need smaller densities.

Solution: a fence of fences



The slots are wider than they are deep, so the total density lies below that of the small fences. (works only for density $< 1/2$)

Iteration yields thinner and thinner and thinner and thinner fences ...

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

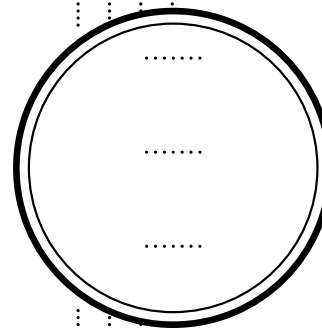
⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮



⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

⋮ ⋮ ⋮ ⋮

Conclusion

We introduced α -Kings (with any $\alpha \in \mathbb{R}^+$) to focus on speed as the essential parameter in the Angel Problem.

Theorem The Devil catches any α -King with $\alpha < 2$.

Conclusion

We introduced α -Kings (with any $\alpha \in \mathbb{R}^+$) to focus on **speed** as the essential parameter in the Angel Problem.

Theorem The Devil catches any α -King with $\alpha < 2$.

Question Can he also catch the 2-King?