# The Density of Iterated Crossing Points and a Gap Result for Triangulations of Finite Point Sets 

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Theorem. For any other set $P$, the limit $P^{\infty}$ is dense in some region of positive measure.

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For any non-stable $P$, the set $P^{\prime \prime \prime \prime}$ contains 5 points in convex position.


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[Ebbers-Baumann et. al., 2005]


## Dilation

$\mathrm{G}=(\mathrm{V}, \mathrm{E})$ plane graph (without crossings).
Define the dilation of G:

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Theorem. [Keil \& Gutwin, 1992]
The dilation of any Dilaunay triangulation is bounded by 2.42.
iterated intersections
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$\Delta<1.02$
[Lorenz, 2004]

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Theorem. Every point set P with $\# \mathrm{P}^{\infty}=\infty$ has dilation $\Delta(\mathrm{P})>1$.

The dilation of any triangulation $T$ that contains $P$ as vertices is bounded away from 1 by some $\gamma_{p}$. There's a "gap!"

## Overview

## Density Theorem.

If $P^{\infty}$ is infinite then it is dense in some region.

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Approximation Lemma. (from exact intersections to dilation >1)

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## Gap Theorem.

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Cor. [Ismailescu \& Radoičić, 2004] If whole lines instead of segments, then the process on 4 pt's in non-convex position densely covers the whole plane.

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## Open Problems:

- What is the dilation of the regular pentagon?
- Is the infimum in $\Delta(\mathrm{P})$ always attained by some triangulation?

