

On Generalizing Decidable Standard Prefix Classes of First-Order Logic

Marco Voigt

Generalized BSR fragment:

relational sentences

$$\forall \bar{x}_1 \exists \bar{y}_1 \dots \forall \bar{x}_n \exists \bar{y}_n. \psi$$

where

the set of atoms in ψ can be partitioned into

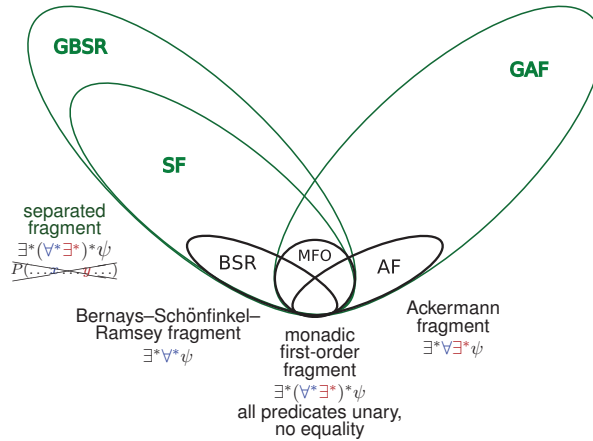
$$\begin{aligned} At_0(\bar{x}_1, \bar{x}_2, \bar{x}_3, \dots, \bar{x}_{n-1}, \bar{x}_n), \\ At_1(\bar{y}_1, \bar{x}_2, \bar{x}_3, \dots, \bar{x}_{n-1}, \bar{x}_n), \\ At_2(\bar{y}_1, \bar{y}_2, \bar{x}_3, \dots, \bar{x}_{n-1}, \bar{x}_n), \\ At_3(\bar{y}_1, \bar{y}_2, \bar{y}_3, \dots, \bar{x}_{n-1}, \bar{x}_n), \\ \vdots \end{aligned}$$

$$\begin{aligned} At_{n-1}(\bar{y}_1, \bar{y}_2, \bar{y}_3, \dots, \bar{y}_{n-1}, \bar{x}_n), \\ At_n(\bar{y}_1, \bar{y}_2, \bar{y}_3, \dots, \bar{y}_{n-1}, \bar{y}_n), \end{aligned}$$

such that

$$\text{vars}(At_i) \cap \text{vars}(At_j) \cap \bar{x}_\ell = \emptyset$$

for all $i \neq j$ and every ℓ .



Generalized Ackermann fragment:

relational sentences

$$\forall \bar{x}_1 \exists \bar{u}_1 \bar{v}_1 \dots \forall \bar{x}_n \exists \bar{u}_n \bar{v}_n. \psi$$

where

the set of atoms in ψ can be partitioned into

$$\begin{aligned} At_{x_1}(\bar{x}_1, \bar{v}_1, \bar{v}_2, \bar{v}_3, \dots, \bar{v}_{n-2}, \bar{v}_{n-1}, \bar{v}_n), \\ At_{x_2}(\bar{u}_1, \bar{x}_2, \bar{v}_2, \bar{v}_3, \dots, \bar{v}_{n-2}, \bar{v}_{n-1}, \bar{v}_n), \\ At_{x_3}(\bar{u}_1, \bar{u}_2, \bar{x}_3, \bar{v}_3, \dots, \bar{v}_{n-2}, \bar{v}_{n-1}, \bar{v}_n), \\ \vdots \end{aligned}$$

$$\begin{aligned} At_{x_{n-1}}(\bar{u}_1, \bar{u}_2, \bar{u}_3, \bar{u}_4, \dots, \bar{x}_{n-1}, \bar{v}_{n-1}, \bar{v}_n), \\ At_{x_n}(\bar{u}_1, \bar{u}_2, \bar{u}_3, \bar{u}_4, \dots, \bar{u}_{n-1}, \bar{x}_n, \bar{v}_n), \\ At_0(\bar{u}_1, \bar{u}_2, \bar{u}_3, \bar{u}_4, \dots, \bar{u}_{n-1}, \bar{u}_n), \end{aligned}$$

such that

$$\text{vars}(At_{x'}) \cap \text{vars}(At_{x''}) \cap \bar{v}_\ell = \emptyset$$

for every $x' \neq x''$ and every ℓ .

Example from $GBSR \cap GAF$: $\exists u \forall x \exists y \forall z. (P(u, z) \wedge \neg Q(u, x, x)) \vee (P(y, u) \wedge Q(u, y, z))$

Computational complexity:

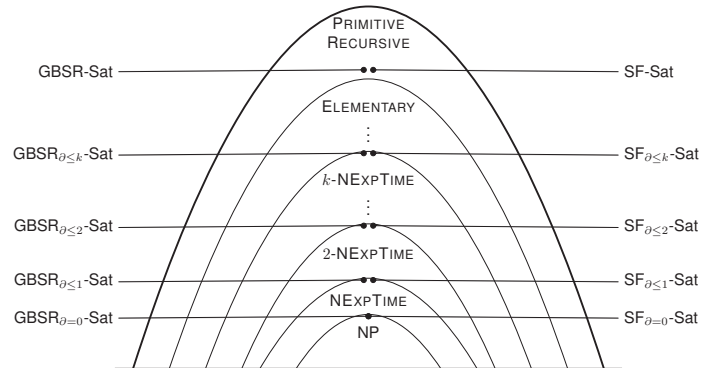
The satisfiability problem for GBSR / SF sentences induces a hierarchy of k -NEXPTIME-complete problems:

The **degree** ∂ of a GBSR sentence measures how many \exists -blocks interact at the level of atoms:

if variables y_1, \dots, y_ℓ from k distinct \exists -blocks are interlinked via atoms

$$P_1(y_1, y_2), P_2(y_2, y_3), \dots, P_n(y_{\ell-1}, y_\ell),$$

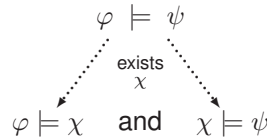
then $\partial \geq k$.



More properties:

- (1) GBSR and GAF possess the **finite model property**:
 - Satisfiable GAF sentences have models over finite domains.
 - Every model of a GBSR sentence contains a finite substructure that is a model, too.
- (2) GBSR and BSR are **equivalent**: Every GBSR sentence can be transformed into an equivalent BSR sentence. In the worst case, this transformation incurs a non-elementary blowup.

- (3) GBSR and SF without equality are **closed under Craig-Lyndon interpolation**:



- φ, ψ are GBSR sentences without equality.
- χ is a BSR sentence over φ 's and ψ 's joint voc.
- Every P with a positive (negative) occurrence in χ occurs positively (negatively) in φ and in ψ .

References

- STURM, VOIGT, WEIDENBACH. *Deciding Satisfiability when Universal and Existential Variables are Separated*. LICS 2016.
 VOIGT. *A Fine-Grained Hierarchy of Hard Problems in the Separated Fragment*. LICS 2017.
 VOIGT. *On Generalizing Decidable Standard Prefix Classes of First-Order Logic*. arXiv:1706.03949 [cs.LO], 2017.

