

**Assignment 1 (DPLL)**

(10 points)

Let  $N$  be the following set of propositional clauses:

$$P \vee Q \vee R \vee S \quad (1)$$

$$P \quad \vee \neg S \vee T \quad (2)$$

$$Q \quad \vee S \quad (3)$$

$$Q \vee \neg R \quad \vee \neg T \quad (4)$$

$$\neg S \vee T \quad (5)$$

$$\neg P \quad \vee R \vee \neg S \vee \neg T \quad (6)$$

$$\neg Q \quad \vee S \vee T \quad (7)$$

$$\neg Q \quad \vee \neg T \quad (8)$$

Use the relation  $\Rightarrow_{\text{DPLL}}$  to test whether  $N$  is satisfiable or not; if it is satisfiable, give a model of  $N$ . Start with the “Decide” rule for the literal  $P$ , then use the “Decide” rule for the literal  $Q$ . If you use the “Backjump” rule, use the *best possible* backjump clause and go to the *best possible* successor state.

**Assignment 2 (Propositional Logic)**

(10 points)

Let  $\Pi$  and  $\Pi'$  be sets of propositional variables and let  $\mu$  be an injective (one-to-one) mapping from  $\Pi$  to  $\Pi'$ . For every propositional formula  $F$  over  $\Pi$ , let  $\mu(F)$  be the formula that one obtains from  $F$  by replacing every propositional variable  $P$  in  $F$  by the propositional variable  $\mu(P)$ . Prove: If  $\mu(F)$  is valid, then  $F$  is valid. (Note: This proof needs an induction argument; write it down in detail.)

**Assignment 3 (Propositional Logic)**

(6 + 6 = 12 points)

**Part (a)**

Prove or refute: If  $F$ ,  $G$ , and  $H$  are propositional formulas,  $\neg F \vee G$  is valid, and  $F \vee H$  is satisfiable, then  $G \vee H$  is satisfiable.

**Part (b)**

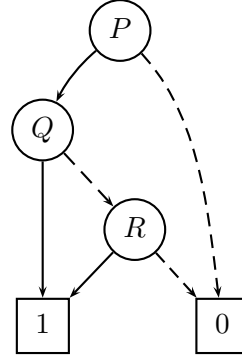
Prove or refute: If  $F$ ,  $G$ , and  $H$  are propositional formulas, and  $(F \wedge H) \rightarrow (G \wedge H)$  is valid, then  $F \rightarrow G$  is valid.

**Assignment 4 (OBDDs)**

(6 + 6 + 6 = 18 points)

**Part (a)**

Give a propositional formula  $F$  that is represented by the reduced OBDD on the right.

**Part (b)**

How many different reduced OBDDs over the propositional variables  $\{P, Q, R\}$  have exactly one interior (non-leaf) node?

**Part (c)**

Find a propositional formula  $G$  over the propositional variables  $\{P, Q, R\}$ , such that the reduced OBDD for  $G$  has three interior nodes and the reduced OBDD for  $F \vee G$  has one interior node. Give the reduced OBDDs for  $G$  and  $F \vee G$ .

**Assignment 5 (Algebras)**

(6 + 6 + 6 = 18 points)

Let  $\Sigma = (\Omega, \Pi)$  with  $\Omega = \{b/0, f/1\}$  and  $\Pi = \{p/1\}$ .

**Part (a)**

How many different Herbrand interpretations over  $\Sigma$  do exist? Explain briefly.

**Part (b)**

How many different Herbrand models over  $\Sigma$  does the following formula  $F$  have?

$$p(b) \wedge \forall x \neg p(f(f(x)))$$

**Part (c)**

Give an example of a  $\Sigma$ -algebra with the universe  $\{1, 2\}$  that is a model of  $F$ .

**Assignment 6 (Termination)**

(12 points)

Let  $\succ$  be a well-founded ordering on the set  $M$ . We define a binary relation  $\triangleright$  on finite subsets of  $M$  in the following way:

$$S \triangleright S \cup \{m_1, \dots, m_k\} \quad \text{if } k \geq 1, \{m_1, \dots, m_k\} \subseteq M, \\ \text{and there exists an } m' \in S \\ \text{such that } m' \text{ is minimal in } S \\ \text{and } m' \succ m_i \text{ for all } i \in \{1, \dots, k\}$$

$$S \triangleright S \setminus \{m'\} \quad \text{if } m' \in S \text{ and } m' \text{ is not minimal in } S$$

Prove that the relation  $\triangleright$  is terminating.