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Tutorials for “Decision Procedures for Logical Theories”  
Exercise sheet 9

**Exercise 9.1:** (7 P.)

Use the superposition calculus to refute the following set of clauses:

$$\begin{aligned} f(g(x)) &\approx 0 \vee g(x) \approx b \\ f(b) &\approx 0 \\ \neg f(g(a)) &\approx 0 \vee g(a) \approx c \\ \neg f(c) &\approx 0 \end{aligned}$$

Use an ordering  $\succ$  that compares two ground terms  $t$  and  $t'$  by lexicographically comparing first the number of  $f$ 's in  $t$  and  $t'$ , then the number of  $g$ 's, then the number of  $a$ 's, then the number of  $b$ 's, and finally the number of  $c$ 's. Use a selection function that selects all trivially false literals (that is, literals of the form  $\neg t \approx t$ ) and nothing else. Perform only inferences that satisfy the conditions of the superposition calculus.

**Exercise 9.2:** (5 P.)

Let  $\mathcal{A}$  and  $\mathcal{A}'$  be  $\Sigma$ -algebras, let  $F$  be a  $\Sigma$ -formula. Prove: If  $\mathcal{A}$  and  $\mathcal{A}'$  are isomorphic, then  $\mathcal{A} \models F$  if and only if  $\mathcal{A}' \models F$ .

**Exercise 9.3:** (3 P.)

Let  $\mathcal{T} = (\mathbb{Q}, +)$ . Use the Nelson/Oppen algorithm  $\mathcal{NO}_{\mathcal{D}}[\mathcal{T}, \Phi]$  to check whether the constraint

$$\begin{aligned} 2x + z &\approx 0 \\ \wedge 2x' + z &\approx 0 \\ \wedge x + y' &\approx y + x' + 1 \\ \wedge f(x, x) &\approx y \\ \wedge f(x, x') &\approx y' \end{aligned}$$

over  $\mathcal{T}^{\Phi}$  is satisfiable.

**Exercise 9.4:** (5 P.)

Prove Theorem 7.1: If all paths in a derivation tree from  $E_1, E_2$  end in  $\perp$ , then  $E_1, E_2$  is unsatisfiable in  $\mathcal{T}_1 + \mathcal{T}_2$ .

Put your solution into the mail box at the door of room 627 in the MPI building (46.1) before January 16, 14:00.