

Tight Bounds for the Approximation Ratio of the Hypervolume Indicator

Karl Bringmann¹ and Tobias Friedrich²

¹ Universität des Saarlandes, Saarbrücken, Germany

² Max-Planck-Institut für Informatik, Saarbrücken, Germany

Abstract The hypervolume indicator is widely used to guide the search and to evaluate the performance of evolutionary multi-objective optimization algorithms. It measures the volume of the dominated portion of the objective space which is considered to give a good approximation of the Pareto front. There is surprisingly little theoretically known about the quality of this approximation. We examine the multiplicative approximation ratio achieved by two-dimensional sets maximizing the hypervolume indicator and prove that it deviates significantly from the optimal approximation ratio. This provable gap is even exponential in the ratio between the largest and the smallest value of the front. We also examine the additive approximation ratio of the hypervolume indicator and prove that it achieves the optimal additive approximation ratio apart from a small factor $\leq n/(n-2)$, where n is the size of the population. Hence the hypervolume indicator can be used to achieve a very good additive but not a good multiplicative approximation of a Pareto front.

1 Introduction

Most real-world optimization problems have to deal with multiple objectives (like time vs. cost) and cannot be easily described by a single objective function. This implies that there is in general no unique optimum, but a possibly very large set of incomparable solutions which forms a Pareto front. Many different multi-objective evolutionary algorithms (MOEAs) have been developed to find a Pareto set of (preferably small) size n which gives a *good approximation* of the Pareto front. A popular way to measure the quality the approximation is the hypervolume indicator. It measures the volume of the dominated space [18]. For a small number of objective, MOEAs which directly use the hypervolume indicator to guide the search are the methods of choice. These include for example the generational MO-CMA-ES [8, 16], the SMS-EMOA [3, 6], and variants of IBEA [17, 19].

One of the reasons why the hypervolume indicator is so popular is that it matches very well with our *intuition* how a good approximation of a Pareto front should look like. However, there is only little known whether maximizing the hypervolume also gives a good approximation of the Pareto front in *strictly mathematical sense*. Considering the wide use of the hypervolume indicator, the

question whether it achieves a good approximation appears to be fundamental. The distribution of the points maximizing the hypervolume indicator has been examined by several authors. It was observed that “convex regions may be preferred to concave regions” [13, 18] as well as that HYP is “biased towards the boundary solutions” [5]. In contrast to this, such sets are empirically “well distributed” according to [6, 9, 10] and it was also proven that for the number of points $n \rightarrow \infty$ the density of points only depends on the gradient [2].

However, the question whether sets maximizing the hypervolume give an approximation of the Pareto front in the mathematical sense remained open besides two preliminary papers [4, 7]. We follow up on this and study the approximation quality of the hypervolume indicator by classic approximation theory. Which concept from approximation theory is the right measure depends on the problem at hand. As a general rule of thumb, for linear axes this is the additive approximation ratio while for exponential axes this is the multiplicative approximation ratio.

To illustrate this with a small example, consider a knapsack problem (see e.g. [18]) with linearly distributed weights and exponentially distributed profits. In this case a good approximation of the front should be an additive approximation of the weights and a multiplicative approximation of the profits. Within this example the result of this paper is that compared to the optimal set with best possible approximation, sets maximizing the hypervolume only achieve the first aim, not necessarily the latter.

In our previous paper [4], we proved that for all possible Pareto fronts the multiplicative approximation factor achieved by a set of n solutions maximizing the hypervolume indicator is $1 + \Theta(1/n)$ (cf. Theorem 3.6)³. As this was shown to be *asymptotically equivalent* to the optimal multiplicative approximation factor (cf. Corollary 3.4), we concluded that the hypervolume indicator is guiding the search in the correct direction for sufficiently large n . However, the size n of a population is usually not large. Also, the constant factors hidden by the Θ might still be larger for the set maximizing hypervolume compared to the set with best possible approximation factor.

Our results

We significantly extend the results of [4]. First, we are now able to give tight bounds on the multiplicative approximation ratio depending on the ratio A/a between the largest and smallest coordinate⁴. Using this notation, the precise result of [4] is the computation of the optimal multiplicative approximation ratio as $1 + \log(A/a)/n$ (cf. Corollary 3.4). We are now able to show that the multiplicative approximation ratio for a set maximizing the hypervolume is strictly

³ The precise statements of this and the following results of this introduction are slightly more technical. For details see the respective theorems.

⁴ The approximation ratio actually depends on the ratios in both dimensions. To simplify the presentation in this introduction, we assume here that the ratio A/a in the first dimension is equal to the ratio B/b in the second dimension.

larger, namely of the order of *at least* $1 + \sqrt{A/a}/n$ (cf. Theorem 3.7). This implies that the dependence of this multiplicative approximation ratio on the ratio A/a can be exponentially worse than in the optimal case. Hence for numerically very wide spread fronts (that is, large A/a) there are Pareto sets which give a much better multiplicative approximation than the Pareto sets which maximize the hypervolume.

Second, we now also analyze the additive approximation ratio of the hypervolume indicator. While the multiplicative approximation factor is determined by the ratio A/a , the additive approximation factor is determined by the width of the domain $A - a$. We prove that the optimal additive approximation ratio is $(A - a)/n$ (cf. Theorem 4.3) and upper bound the additive approximation ratio achieved by a set maximizing the hypervolume by $(A - a)/(n - 2)$ (cf. Theorem 4.5). This is a very strong statement, as apart from the small factor $n/(n - 2)$ the additive approximation ratio achieved when maximizing the hypervolume is optimal! This shows that the hypervolume indicator yields a much *better additive than multiplicative* approximation.

2 Preliminaries

We only consider the case of two objectives where there is a mapping from an arbitrary search space to an objective space which is a subset of \mathbb{R}^2 . Throughout this paper, we will only work on the objective space. For points from the objective space we define the following dominance relation:

$$\begin{aligned} (x_1, y_1) \preceq (x_2, y_2) &\text{ iff } x_1 \leq x_2 \text{ and } y_1 \leq y_2, \\ (x_1, y_1) \prec (x_2, y_2) &\text{ iff } (x_1, y_1) \preceq (x_2, y_2) \text{ and } (x_1, y_1) \neq (x_2, y_2). \end{aligned}$$

We restrict ourselves to Pareto fronts that can be written as $\{(x, f(x)) \mid x \in [a, A]\}$ where $f: [a, A] \rightarrow [b, B]$ is a (not necessarily strictly) monotonically decreasing, upper semi-continuous⁵ function with $f(a) = B$, $f(A) = b$ for $a < A$, $b < B$. We write $\mathcal{F} = \mathcal{F}_{[a,A] \rightarrow [b,B]}$ for the set of all such functions f . We will use the term *front* for both, the set of points $\{(x, f(x)) \mid x \in [a, A]\}$, and the function f .

The condition of f being upper semi-continuous cannot be relaxed further as without it the f lacks symmetry in the two objectives in the following sense: Being upper semi-continuous is necessary and sufficient for the existence of the inverse function $f^{-1}: [b, B] \rightarrow [a, A]$ defined by setting $f^{-1}(y) := \max\{x \in [a, A] \mid f(x) \geq y\}$. Without upper semi-continuity this maximum does not exist in general. Furthermore, this condition implies that there is a set maximizing the hypervolume indicator.

⁵ Semi-continuity is a weaker property than normal continuity. A function f is said to be upper semi-continuous if for all points x of its domain, $\limsup_{y \rightarrow x} f(y) \leq f(x)$. Intuitively speaking this means that for all points x the function values for arguments near x are either close to $f(x)$ or less than $f(x)$. For more details see e.g. Rudin [15].

Note that the set \mathcal{F} of fronts we consider is a very general one. Most papers that theoretically examine the hypervolume indicator assume that the front is continuous and differentiable (e.g. [1, 2, 7]), and are thus not able to give results about discrete fronts, which we can.

Let $n \in \mathbb{N}$. For fixed $[a, A], [b, B] \subset \mathbb{R}$ we call a set $P = \{p_1, \dots, p_n\} \subset [a, A] \times [b, B]$ a *solution set* (of size n) and write $\mathcal{P} = \mathcal{P}_n$ for the set of all such solution sets. A solution set P is said to be *feasible* for a front $f \in \mathcal{F}$, if $y \leq f(x)$ for all $p = (x, y) \in P$. We write $\mathcal{P}^f \subseteq \mathcal{P}$ for the set of all solution sets that are feasible for f .

We are now ready to formally define the hypervolume indicator. It was first introduced for performance assessment in multiobjective optimization by Zitzler and Thiele [18], but since then also has become a very popular way to guide the search in multi-objective evolutionary optimizers. On a two-dimensional objective space it is defined as follows.

Definition 2.1 *The hypervolume indicator $\text{HYP}(P)$ of a solution set $P \in \mathcal{P}$ relative to a reference point $R = (R_x, R_y)$ is*

$$\text{HYP}(P) := \text{VOL} \left(\bigcup_{(x,y) \in P} [R_x, x] \times [R_y, y] \right).$$

with $\text{VOL}(\cdot)$ being the usual Lebesgue measure.

3 Multiplicative Approximation

The standard measure of approximation quality in approximation theory is the multiplicative approximation ratio. We use the multi-objective definition for the multiplicative approximation ratio by Papadimitriou and Yannakakis [14] which was also used in [4, 7, 11, 12]. Note that here and in the rest of the paper when talking about multiplicative approximation we require $a, b > 0$ as this ratio only makes sense for positive values.

Definition 3.1 *Let $f \in \mathcal{F}$ and $P \in \mathcal{P}^f$. The solution set P is a multiplicative α -approximation of f if for each $\hat{x} \in [a, A]$ there is a point $p = (x, y) \in P$ with*

$$\hat{x} \leq \alpha x \quad \text{and} \quad f(\hat{x}) \leq \alpha y$$

where $\alpha \in \mathbb{R}$, $\alpha \geq 1$. The multiplicative approximation ratio of P with respect to f is then defined as

$$\alpha^*(f, P) := \inf \{ \alpha \in \mathbb{R} \mid P \text{ is a multiplicative } \alpha\text{-approximation of } f \}.$$

The quality of an algorithm which calculates a solution set of size n for each Pareto front in \mathcal{F} has to be compared with the respective optimal approximation ratio defined as follows.

Definition 3.2 For fixed $[a, A]$, $[b, B]$, and n , let

$$\alpha_{OPT}^* := \sup_{f \in \mathcal{F}} \inf_{P \in \mathcal{P}^f} \alpha^*(f, P).$$

The value α_{OPT}^* is chosen such that every front in \mathcal{F} can be approximated by n points to a ratio of α_{OPT}^* , and there is a front which cannot be approximated better. In [4] the authors showed the following two results.

Theorem 3.3 (from [4]) $\alpha_{OPT}^* = \min\{A/a, B/b\}^{1/n}$.

Corollary 3.4 (from [4]) For all $n \geq \log(\min\{A/a, B/b\})/\varepsilon$ and $\varepsilon \in (0, 1)$,

$$\begin{aligned} \alpha_{OPT}^* &\geq 1 + \frac{\log(\min\{A/a, B/b\})}{n}, \\ \alpha_{OPT}^* &\leq 1 + (1 + \varepsilon) \frac{\log(\min\{A/a, B/b\})}{n}. \end{aligned}$$

We further want to measure the approximation of the solution set of size n maximizing HYP. As there might be several solution sets maximizing HYP, we consider the worst case and use the following definition.

Definition 3.5 For fixed $[a, A]$, $[b, B]$, and n , let

$$\mathcal{P}_{HYP}^f := \{P \in \mathcal{P}^f \mid \text{HYP}(P) = \max_{Q \in \mathcal{P}^f} \text{HYP}(Q)\} \text{ for } f \in \mathcal{F}, \text{ and}$$

$$\alpha_{HYP}^* := \sup_{f \in \mathcal{F}} \sup_{P \in \mathcal{P}_{HYP}^f} \alpha^*(f, P).$$

The set \mathcal{P}_{HYP}^f is the set of all feasible solution sets that maximize HYP on f . The value α_{HYP}^* is chosen such that for every front f in \mathcal{F} every solution set maximizing HYP approximates f by a ratio of at most α_{HYP}^* . Note that this assumes that there is at least one solution set which maximizes the indicator, i.e., the set \mathcal{P}_{HYP}^f is non-empty. That this is indeed the case was proven in [4].

In [4] the authors also examined α_{HYP}^* and showed an upper bound that has the same asymptotic behavior as α_{OPT}^* , but a much larger constant factor.

Theorem 3.6 (from [4]) Let $f \in \mathcal{F}$, $n > 4$, and let $R = (R_x, R_y) \preceq (0, 0)$ be the reference point. If we have

- $n \geq 2 + \max\{\sqrt{A/a}, \sqrt{B/b}\}$ or
- $R_x \leq -\sqrt{Aa}/n$, $R_y \leq -\sqrt{Bb}/n$,

then

$$\alpha_{HYP}^* \leq 1 + \frac{\sqrt{A/a} + \sqrt{B/b}}{n - 4}.$$

The previous paper [4] left open whether (i) the upper bound of Theorem 3.6 is not tight and α_{HYP}^* is actually much closer to the bounds for α_{OPT}^* given in Corollary 3.4 or (ii) α_{HYP}^* is indeed significantly larger than α_{OPT}^* . By giving a lower bound for α_{HYP}^* we can now prove the latter. In the following theorem we restrict ourselves to the case of $A/a = B/b$. We show that in this situation the bound of Theorem 3.6 is tight except for a small constant factor.

Theorem 3.7 *Let $n \geq 4$, $A/a = B/b \geq 13$, and $R = (R_x, R_y) \preceq (0, 0)$ be the reference point. Then*

$$\alpha_{HYP}^* \geq 1 + \frac{2\sqrt{A/a - 1}}{3(n-1)}.$$

The proof of this theorem will be provided in the full version of the paper.

4 Additive Approximation

After the previous section showed that sets maximizing the hypervolume have sub-optimal multiplicative approximation ratio we now analyze their additive approximation properties. Analogous to Definition 3.1 we use the following definition.

Definition 4.1 *Let $f \in \mathcal{F}$ and $P \in \mathcal{P}^f$. The solution set P is an additive α -approximation of f if for each $\hat{x} \in [a, A]$ there is a point $p = (x, y) \in P$ with*

$$\hat{x} \leq x + \alpha \quad \text{and} \quad f(\hat{x}) \leq y + \alpha$$

where $\alpha \in \mathbb{R}$, $\alpha \geq 0$. The additive approximation ratio of P with respect to f is defined as

$$\alpha^+(f, P) := \inf\{\alpha \in \mathbb{R} \mid P \text{ is an additive } \alpha\text{-approximation of } f\}.$$

Again, we are interested in the optimal approximation ratio for Pareto fronts in \mathcal{F} . Analogous to Definition 3.2 we give the following definition.

Definition 4.2 *For fixed $[a, A]$, $[b, B]$, and n , let*

$$\alpha_{OPT}^+ := \sup_{f \in \mathcal{F}} \inf_{P \in \mathcal{P}^f} \alpha^+(f, P).$$

Analogously to the precise bound $\alpha_{OPT}^* = \min\{A/a, B/b\}^{1/n}$ of Theorem 3.3 for the optimal multiplicative approximation ratio, we can prove the following for the optimal additive approximation ratio α_{OPT}^+ .

Theorem 4.3 $\alpha_{OPT}^+ = \frac{\min\{A - a, B - b\}}{n}$.

The proof of Theorem 4.3 will be provided in the full version of the paper.

In order to compare the optimal additive approximation ratio with the approximation ratio achieved by the hypervolume, we give the following definition analogously to the definition of α_{HYP}^* in Definition 3.5.

Definition 4.4 *For fixed $[a, A]$, $[b, B]$, n , and $f \in \mathcal{F}$ let*

$$\alpha_{HYP}^+ := \sup_{f \in \mathcal{F}} \sup_{P \in \mathcal{P}_{HYP}^f} \alpha^+(f, P).$$

We can now state the main result of this paper that α_{HYP}^+ is very close to α_{OPT}^+ . Similar to the proof of the upper bound for α_{HYP}^* of Theorem 3.6 we can prove the following upper bound for α_{HYP}^+ .

Theorem 4.5 *For all $n > 2$ and $(n - 2) \min\{a - R_x, b - R_y\} \geq \sqrt{(A - a)(B - b)}$,*

$$\alpha_{HYP}^+ \leq \frac{\sqrt{(A - a)(B - b)}}{n - 2}.$$

Let us briefly discuss the result before the theorem will be proven in the remainder of this section. First note that the precondition is fulfilled if n is large enough or if the reference point is sufficiently far away from (a, b) . Hence this is no real restriction. Moreover, compare this result to the bound for the optimal additive approximation ratio of Theorem 4.3. This shows that for $A - a \approx B - b$ and moderately sized n , α_{HYP}^+ is very close to α_{OPT}^+ . More precisely, for $A - a \ll B - b$ (or $A - a \gg B - b$) the constant in Theorem 4.5 is the geometric mean of $A - a$ and $B - b$ while in Theorem 4.3 it is instead the minimum of both. As there is a provable gap of log vs. square root of A/a for the multiplicative approximation ratio, this proves that HYP yields a much better additive approximation than a multiplicative one.

Proof of Theorem 4.5. Let P be a solution set maximizing HYP on a front $f \in \mathcal{F}$, i.e., $P \in \mathcal{P}_{HYP}^f$. Assume that there are points $p, q \in P$ with $p \prec q$. Such a “redundant” set can maximize HYP only on degenerate fronts: If there is a point $r = (x, f(x))$ on the front which is not dominated by any point in P , then⁶ $P' := P + r - p$ would have $\text{HYP}(P') > \text{HYP}(P)$, as it dominates all the space P dominates united with the space r dominates. Thus, there is no such point r and P dominates already the whole front. In this case the approximation ratio $\alpha^+(f, P) = 1$ and the inequality we want to show holds trivially. This can only happen for f being a step function with less than n steps.

Hence, for the rest of the proof we can assume that there are no points $p, q \in P$ with $p \prec q$. Then we can write $P = \{p_1, \dots, p_n\}$, $p_i = (x_i, y_i)$ with $a \leq x_1 < \dots < x_n \leq A$ and $B \geq y_1 > \dots > y_n \geq b$. Furthermore, we can assume that $y_i = f(x_i)$ as otherwise $P - p_i + p'_i$ with $p'_i = (x_i, f(x_i))$ would have a larger hypervolume than P (this uses that the points in P are mutually non-dominating).

We want to argue about the contribution of a point p to the hypervolume of a solution set P , namely $\text{CON}_P(p) := \text{HYP}(P) - \text{HYP}(P - p)$. In particular we need the minimal contribution of any of the points p_2, \dots, p_{n-1} :

$$\begin{aligned} \text{MINCON}(P) &:= \min_{1 < i < n} \text{CON}_P(p_i) \\ &= \min_{1 < i < n} (x_i - x_{i-1})(f(x_i) - f(x_{i+1})). \end{aligned}$$

⁶ To increase the readability, for a set $P \subset \mathbb{R}^2$ and a point $r \in \mathbb{R}^2$ we define $P + r := P \cup \{r\}$ and $P - r := P \setminus \{r\}$ here and in the remainder of this section.

This value has been (with slightly different notation) examined in [4]. In particular, the authors showed that for $n > 2$

$$\text{MINCON}(P) \leq \frac{(x_n - x_1)(f(x_1) - f(x_n))}{(n - 2)^2}.$$

This implies

$$\text{MINCON}(P) \leq \frac{(A - a)(B - b)}{(n - 2)^2}. \quad (1)$$

Let $r = (x, f(x))$, $x \in [a, A]$ be an arbitrary point and let $\alpha > 0$ be such that r is not additively approximated by α . We make a case distinction depending on the position of r . Let us first assume that r is an “inner point”, i.e., there is an $i \in \{1, \dots, n - 1\}$ with $x_i \leq x < x_{i+1}$. As r is not additively approximated by α , we have

$$x > x_i + \alpha \quad \text{and} \quad f(x) > f(x_{i+1}) + \alpha. \quad (2)$$

As P maximizes the hypervolume indicator on f , replacing the point $p \in P$ contributing $\text{MINCON}(P)$ to P by the point r must not increase the hypervolume. Therefore,

$$\begin{aligned} \text{HYP}(P) &\geq \text{HYP}(P + r - p) = \text{HYP}(P) - \text{CON}_P(p) + \text{CON}_{P+r-p}(r) \\ &\geq \text{HYP}(P) - \text{CON}_P(p) + \text{CON}_{P+r}(r), \end{aligned}$$

which in turn implies

$$\text{MINCON}(P) = \text{CON}_P(p) \geq \text{CON}_{P+r}(r) = (x - x_i)(f(x) - f(x_{i+1})) \stackrel{(2)}{>} \alpha^2.$$

Using equation (1) and taking square roots on both sides gives the desired

$$\alpha < \frac{\sqrt{(A - a)(B - b)}}{n - 2}.$$

It remains to study the case where $r = (x, f(x))$ is an “outer point” with $x \leq x_1$ or $x \geq x_n$. It suffices to examine $x \leq x_1$ as then the case $x \geq x_n$ follows by symmetry in the two objectives.

As r is not approximated by a ratio of α we have $f(x) > f(x_1) + \alpha$. Additionally, replacing the point $p \in P$ contributing $\text{MINCON}(P)$ to P by r must not increase the hypervolume, so we have

$$\begin{aligned} \text{MINCON}(P) &\geq \text{CON}_{P+r-p}(r) \geq \text{CON}_{P+r}(r) = (a - R_x)(f(x) - f(x_1)) \\ &\geq (a - R_x)\alpha. \end{aligned}$$

We use equation (1) again and get

$$\alpha \leq \frac{(A - a)(B - b)}{(a - R_x)(n - 2)^2} \leq \sqrt{(A - a)(B - b)} / (n - 2),$$

where the second inequality follows from the precondition of the theorem. \square

	Multiplicative approximation	Additive approximation
OPT	$1 + \frac{\log(\min\{A/a, B/b\})}{n}$ (Cor. 3.4)	$\frac{\min\{A-a, B-b\}}{n}$ (Thm. 4.3)
HYP	$1 + \frac{\sqrt{A/a} + \sqrt{B/b}}{n-4}$ (Thm. 3.6)	$\frac{\sqrt{(A-a)(B-b)}}{n-2}$ (Thm. 4.5)

Table 1: Results for the optimal approximation ratio and upper bounds for the approximation ratios of HYP. See the cited theorems for the precise statements.

5 Conclusion

Many modern MOEA use the hypervolume indicator to guide the search process. We presented a mathematically rigorous framework to analyze the approximation ratio achieved by sets maximizing the hypervolume. We prove that sets maximizing HYP do *not* give a perfect multiplicative approximation. The proven bounds can be found in Table 1. The multiplicative approximation ratio of HYP is getting large for numerically wide spread fronts with large A/a . On the other hand, we can prove that maximizing HYP gives a close-to-optimal additive approximation.

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