



Michael Sagraloff

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Computer Algebra

<https://www.mpi-inf.mpg.de/departments/algorithms-complexity/teaching/winter17/comp-alg/>

Assignment sheet 9

due: Monday, January 8

Exercise 1: Bounds on the intermediate values in the EEA (4 points)

We consider the Extended Euclidean Algorithm for integer polynomials $f, g \in \mathbb{Z}[x]$ of degree bounded by n and coefficients of absolute value less than 2^τ as presented in the lecture. As usual, define $s_i, t_i, \rho_i, r_i, q_i$ as

$$\begin{aligned} \rho_0 &:= \text{LC}(f), & r_0 &:= \text{normal}(f), & s_0 &:= \rho_0^{-1}, & t_0 &:= 0, \\ \rho_1 &:= \text{LC}(g), & r_1 &:= \text{normal}(g), & s_1 &:= 0, & t_1 &:= \rho_1^{-1} \end{aligned}$$

and, for $1 \leq i \leq \ell$ (with ℓ the index such that $r_\ell \neq 0$ and $r_{\ell+1} = 0$),

$$\begin{aligned} q_i &:= r_{i-1} \text{ quo } r_i, & \rho_{i+1} &:= \text{LC}(r_{i-1} \text{ rem } r_i), & r_{i+1} &:= \text{normal}(r_{i-1} \text{ rem } r_i), \\ s_{i+1} &:= (s_{i-1} - q_i s_i) / \rho_{i+1}, & t_{i+1} &:= (t_{i-1} - q_i t_i) / \rho_{i+1}. \end{aligned}$$

Show that there exist integers μ_i of length $O(n(\tau + \log n))$ such that $\mu_i \cdot \rho_i$ and $\mu_i \cdot q_i$ are integers (integer polynomials) of length (with coefficients of length) $O(n(\tau + \log n))$!

Proceed as follows:

1. Use that a comparable result has already been shown for s_i, t_i , and r_i !
2. Recall that

$$\begin{aligned} R_i &= \begin{pmatrix} s_i & t_i \\ s_{i+1} & t_{i+1} \end{pmatrix} = R_0 \cdot \prod_{j=1}^i Q_j, \quad \text{where} \\ R_0 &= \begin{pmatrix} s_0 & t_0 \\ s_1 & t_1 \end{pmatrix} \quad \text{and} \quad Q_j = \begin{pmatrix} 0 & 1 \\ \rho_{j+1}^{-1} & -q_j \rho_{j+1}^{-1} \end{pmatrix} \end{aligned}$$

and, in particular, $\begin{vmatrix} s_i & t_i \\ s_{i+1} & t_{i+1} \end{vmatrix} = (-1)^{i-1} (\rho_0 \cdots \rho_i)^{-1}$. Use these identities to derive a bound on the length of the numerator and denominator of ρ_i .

3. Prove that $f = q \cdot g$ with $f, g \in \mathbb{Z}[x]$ and $q \in \mathbb{Q}[x]$ implies that there exists a $\lambda \in \mathbb{Z}$ with $|\lambda| < 2^\tau$ such that

$$\lambda \cdot q \in \mathbb{Z}[x] \quad \text{and} \quad \|\lambda q\|_\infty = 2^{O(n+\tau)}.$$

4. Use the fact that $r_{i-1} = q_i r_i + \rho_{i+1} r_{i+1}$ and the previous result to derive a bound on the size of q_i .

Exercise 2: Modular GCD computation (4 points)

Let $f, g \in \mathbb{Z}[x]$ be integer polynomials of degree bounded by n and coefficients of absolute value less than 2^τ , let p be prime such that $p \nmid \text{LC}(f)$ and $p \nmid \text{LC}(g)$, and define $d := \deg \gcd(f, g)$ to be the degree of the GCD of f and g .

1. Show that

$$\gcd(f, g) \equiv \gcd(\bar{f}, \bar{g}) \pmod{p} \quad \text{if and only if} \quad p \nmid \text{sres}_d(f, g),$$

where \bar{f} and \bar{g} are the modular images of f and g in $\mathbb{Z}/p\mathbb{Z}[x]$.

2. Develop a modular algorithm to compute *under guarantee* the degree d of $\gcd(f, g) \in \mathbb{Z}[x]$ and determine its bit complexity in terms of n and τ .

Exercise 3: Applications of subresultants (4 points)

- (a) Let

$$\begin{aligned} f &= x^3 + 4x^2 - 2ax - a^2 \quad \text{and} \\ g &= x^2 - 2a^2. \end{aligned}$$

Choose a such that $\deg \gcd(f, g) = 1$.

- (b) Determine the gcd of

$$\begin{aligned} f &= x^2 + \left(\frac{1}{10}\sqrt{5} - \frac{3}{10}\right)x + \left(\frac{3}{50}\sqrt{5} - \frac{7}{50}\right) \quad \text{and} \\ g &= 4x^2 + \left(-\frac{1}{10}\sqrt{5} + \frac{3}{10}\right)x + \left(\frac{1}{25}\sqrt{5} - \frac{4}{25}\right). \end{aligned}$$

Exercise 4: Errors in the Script (many points)

Detect errors in the script! Depending on the error you will receive the following amount of points for each error:

- Typo: 0.25 point
- Any kind of language error (verified by a native speaker): 0.25 point
- Error in a calculation: 0.5 point
- Serious error (bad argument, etc.): 2 points

In addition, any reasonable recommendation regarding the presentation (explanation, style, etc.) is rewarded with 0.5 point.